

Common Gnoseological Meaning of Goedel and Caratheodory Theorems

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1. Introduction

Vienna Circle

logical positivism - physicalism

Rudolph Carnap Otto von Neurath :

1931 - 1935

*"Any scientifically meaningful statement
is expressible in physical terms*

*- about a movement in the observable space and time -
or, if the statement is not expressible this way*

it is meaningful scientifically

*when it is convertible to a statement about a language,
otherwise it is of no scientific meaning."*

⇒

mutual **relation** among **structures/languages** of:

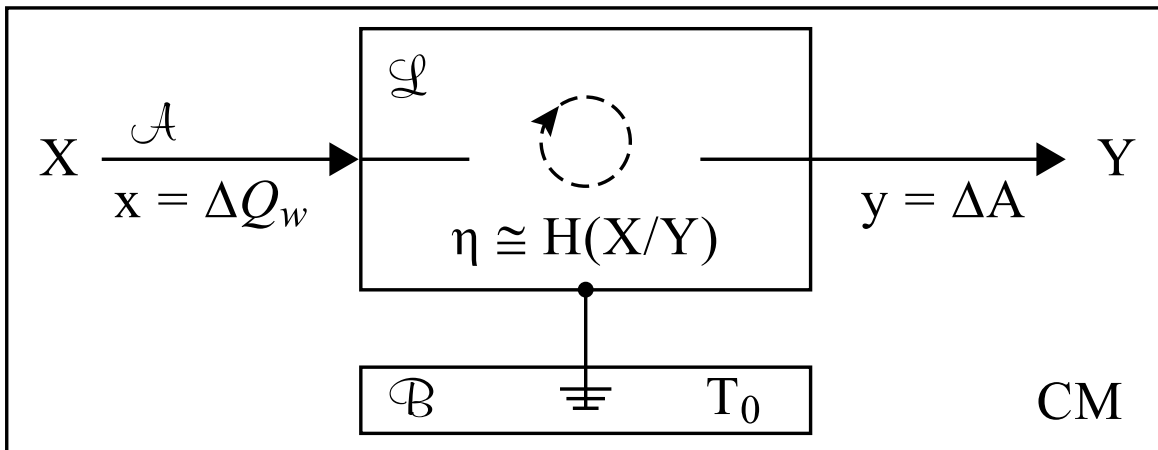
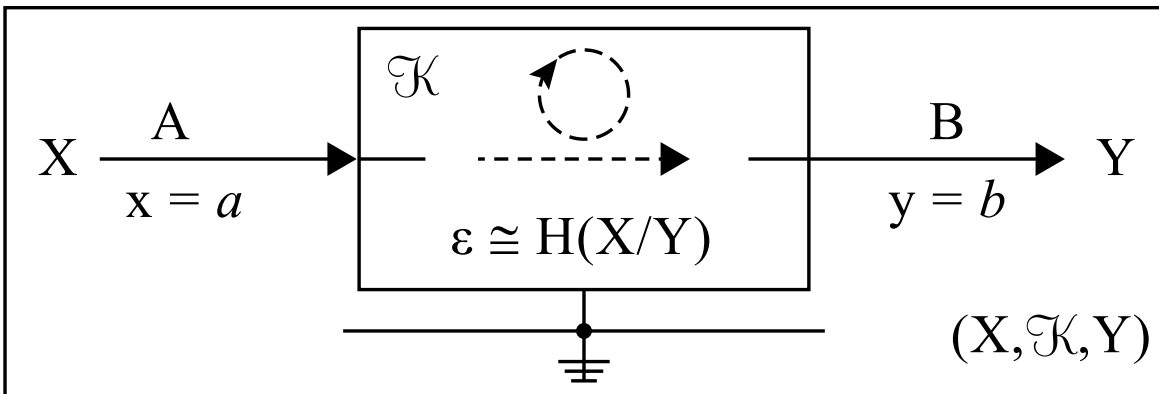
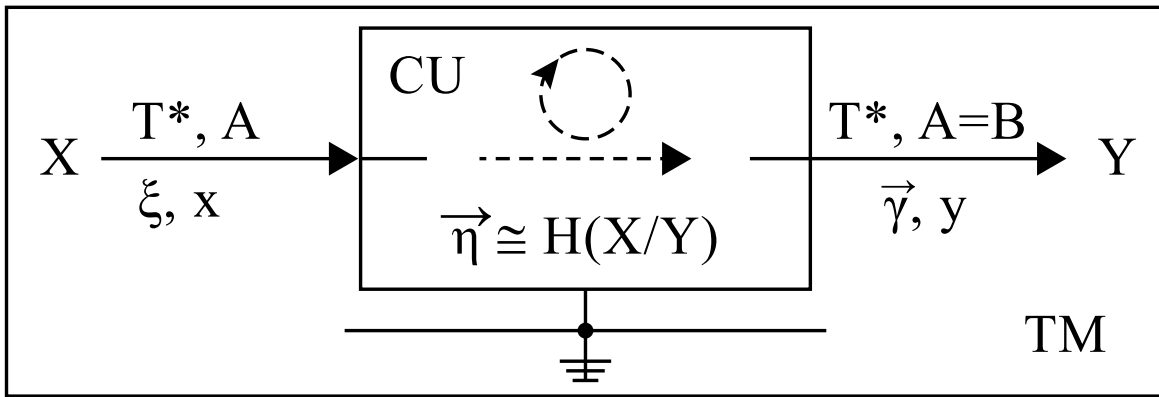
Thermodynamics – **Energy Transformation**

Information Theory – **Message/Information Transfer**

Computing Theory – **Computing/Inference**

Adiabatic Theory

<p style="text-align: center;">COMPUTING THEORY</p>	<p style="text-align: center;">TURING MACHINE <i>TM</i> FINITE-STATE CONTROL UNIT</p>	<p style="text-align: center;">COMPUTING/INFERENCE PROCESS</p> <p style="text-align: center;">STATE/CONFIGURATION TRANSFORMATIONS</p>
<p style="text-align: center;">INFORMATION THEORY</p>	<p style="text-align: center;">SHANNON TRANSFER CHAIN (X, \mathcal{K}, Y) TRANSFER CHANNEL</p>	<p style="text-align: center;">MESSAGE/INFORMATION TRANSFER PROCESS</p>
<p style="text-align: center;">THERMODYNAMICS</p>	<p style="text-align: center;">CARNOT MACHINE <i>CM</i> CARNOT CYCLE</p>	<p style="text-align: center;">HEAT ENERGY TRANSFORMATION PROCESS</p>
<p style="text-align: center;">THERMODYNAMIC ADIABATIC THEORY</p>	<p style="text-align: center;">THERMODYNAMIC ADIABATIC SYSTEM \mathcal{L}</p>	<p style="text-align: center;">STATES' DEVELOPING PROCESS</p>



Processes in all these structures run in the finite physical world and follow its laws.

We model them by the states' $\theta_{[\cdot]}^{\mathcal{L}}$ trajectories $l_{\Omega_{\mathcal{L}}}$ within the heat isolated

$$- \underbrace{[d]Q_{\text{Ext}}=0}_{\text{adiabatic system } \mathcal{L}/\Omega_{\mathcal{L}}} -$$

where is valid:

Caratheodory common formulation of the II. P.T. :

In the arbitrary vicinity of every state of the state space $\Omega_{\mathcal{L}}$ of the adiabatic system \mathcal{L} exist states not reachable from the starting state adiabatically ($[d]Q_{\text{Ext}}=0$) (or the states not reachable by the system at all).

1

For the consistency of the **Peano arithmetic theory** $\mathcal{T}_{\mathcal{P}_A}$ the analog is expressed by:

Gödel incompleteness theorems:²

For the theory $\mathcal{T}_{\mathcal{P}_A}$ exists the true ("1") CLAIM that either this CLAIM and its NEGATION is NOT PROVABLE within the system $\mathcal{P}/\mathcal{T}_{\mathcal{P}_A}$.

3

- CLAIM about the $\mathcal{T}_{\mathcal{P}_A}$ consistency especially -

The CLAIM saying that theory $\mathcal{T}_{\mathcal{P}_A}$ is consistent is not PROVABLE by its means (\mathcal{P}) - by itself.

4

¹Along the given the trajectory $l_{\Omega_{\mathcal{L}}}$ with the given starting point, reversibly or irreversibly. Or such states which are the part of the $\mathcal{L}/\Omega_{\mathcal{L}}$'s *outer construction* and thus of the whole $l_{\Omega_{\mathcal{L}}}$'s *definition*.

²Rosser-Gödel theorem.

³Far from (!) "In..." Attempts to prove/*TO PROVE*/*INFER* it within the system $\mathcal{P}/\mathcal{T}_{\mathcal{P}_A}$ leads to the inconsistency of the consistent (!) system \mathcal{P}_{κ} (in fact we are entering into the inconsistent metasystem \mathcal{P}^{\star} - the real sense of the Proposition V).

⁴It is the *META-CLAIM* not writable within the $\mathcal{T}_{\mathcal{P}_A}$ language.

The adiabatic trajectories $l_{\Omega_{\mathcal{L}}}$ - within the $\mathcal{L}/\Omega_{\mathcal{L}}$ ⁵

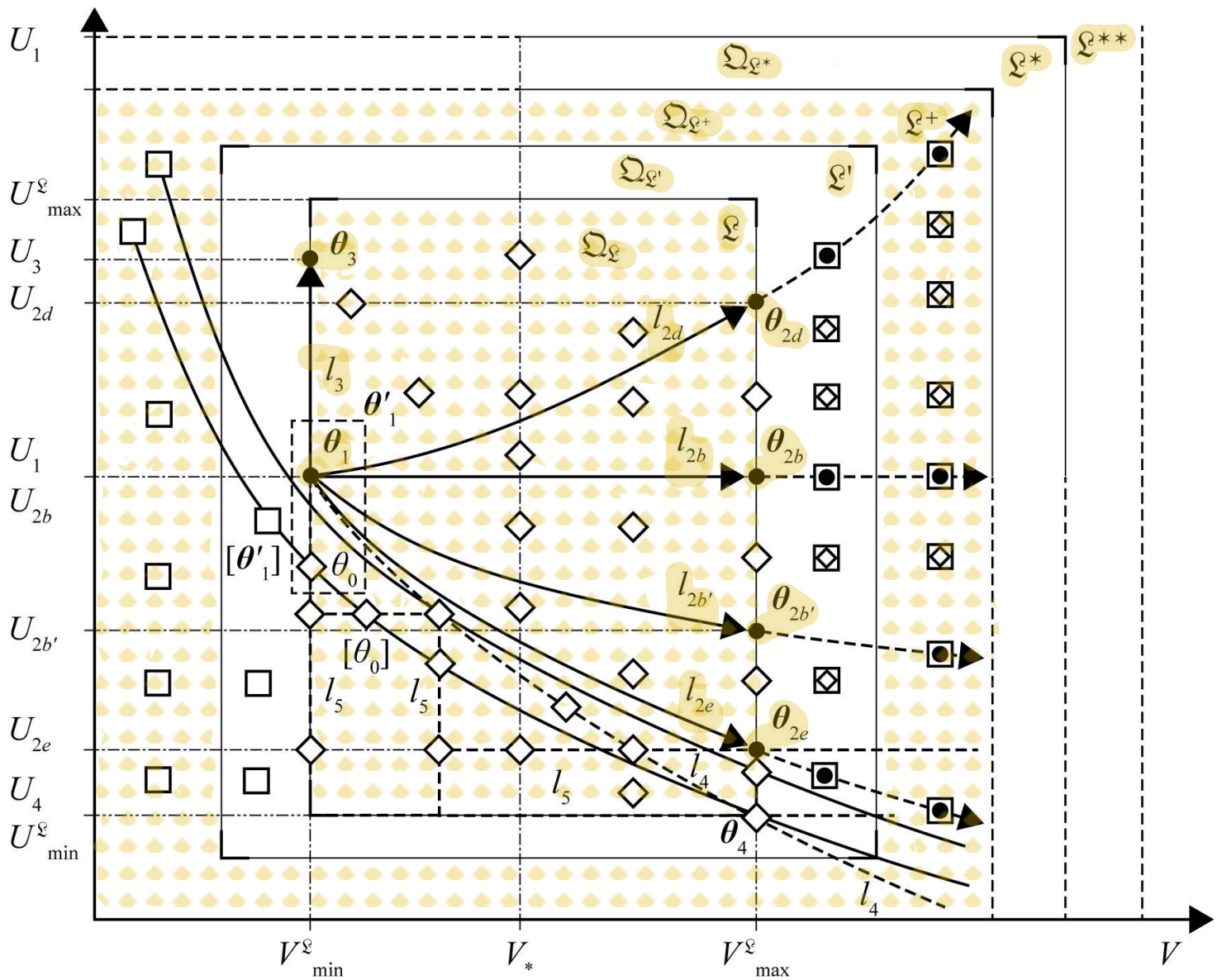
l_{2b} *isothermic irreversible expansion*

$l_{2b'}$ *adiabatic irreversible expansion*

l_{2d} *izobaric irreversible expansion*

l_{2e} *izentropic reversible expansion*

l_3 *izochoric irreversible change*



l_4 *not possible*

⁵Which is the part of the wider meta-language with the vocabulary $\{p, V, T\}$ in which we construct/define this adiabatic state space $\Omega_{\mathcal{L}}$ /system \mathcal{L} by its complement $\{p, V, T\} - \Omega_{\mathcal{L}}$ and the trajectory $l_{\Omega_{\mathcal{L}}}$ by its complement $\Omega_{\mathcal{L}} - l_{\Omega_{\mathcal{L}}}$ $[\{p, V, T\} - (\Omega_{\mathcal{L}} - l_{\Omega_{\mathcal{L}}})]$ but, including its initial state from $\Omega_{\mathcal{L}}$ - from outside.

Peano Axioms/Inference system \mathcal{P} /Theory $\mathcal{T}_{\mathcal{P}A}$

$$1/\mathcal{P} \quad \mathbb{N}_0 = \mathbb{N} \cup \{\mathbf{0}\};$$

$$2 \quad \forall_{\mathbf{x} \in \mathbb{N}_0} | [\exists_{\mathbf{y} \in \mathbb{N}} | [\mathbf{y} = \mathbf{f}(\mathbf{x})]];$$

$$3/\mathcal{P} \quad \forall_{\mathbf{x} \in \mathbb{N}_0} | [\mathbf{0} \neq \mathbf{f}(\mathbf{x})];$$

$$4 \quad \forall_{\mathbf{x} \in \mathbb{N}_0} | [[\mathbf{f}(\mathbf{x}) \neq \mathbf{f}(\mathbf{y})] \Rightarrow (\mathbf{x} \neq \mathbf{y})];$$

5/ \mathcal{P} *axiom/axiomatic schema of the mathematical induction:*

$$[[\varphi(\mathbf{0}) \wedge \forall_{\mathbf{x} \in \mathbb{N}_0} | \varphi(\mathbf{x}) \Rightarrow \varphi[\mathbf{f}(\mathbf{x})]] \Rightarrow \forall_{\mathbf{x} \in \mathbb{N}_0} | \varphi(\mathbf{x})]$$

Inference rule *Modus Ponens*^a

$$\frac{\vdash \mathbf{b}, \vdash (\mathbf{b} \Rightarrow \mathbf{c})}{\vdash \mathbf{c}}, \quad \mathbf{c} - \textit{immediate consequence of } \mathbf{b}$$

^aBesides **the Generalization**. The **Substitution function** from the *Principia Mathematica* is used for the evaluation of the variables be their values or their quantification.

”1” - arithmeticity of the \mathcal{P}

\cong

adiabaticity of the $\mathcal{L}/\Omega_{\mathcal{L}}$.

Consistent $\mathcal{T}_{\mathcal{P}A}$ inference within \mathcal{P}

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moving along trajectories $l_{\Omega_{\mathcal{L}}}$ within the $\Omega_{\mathcal{L}}/\mathcal{L}$.

The states on the adiabatic trajectories

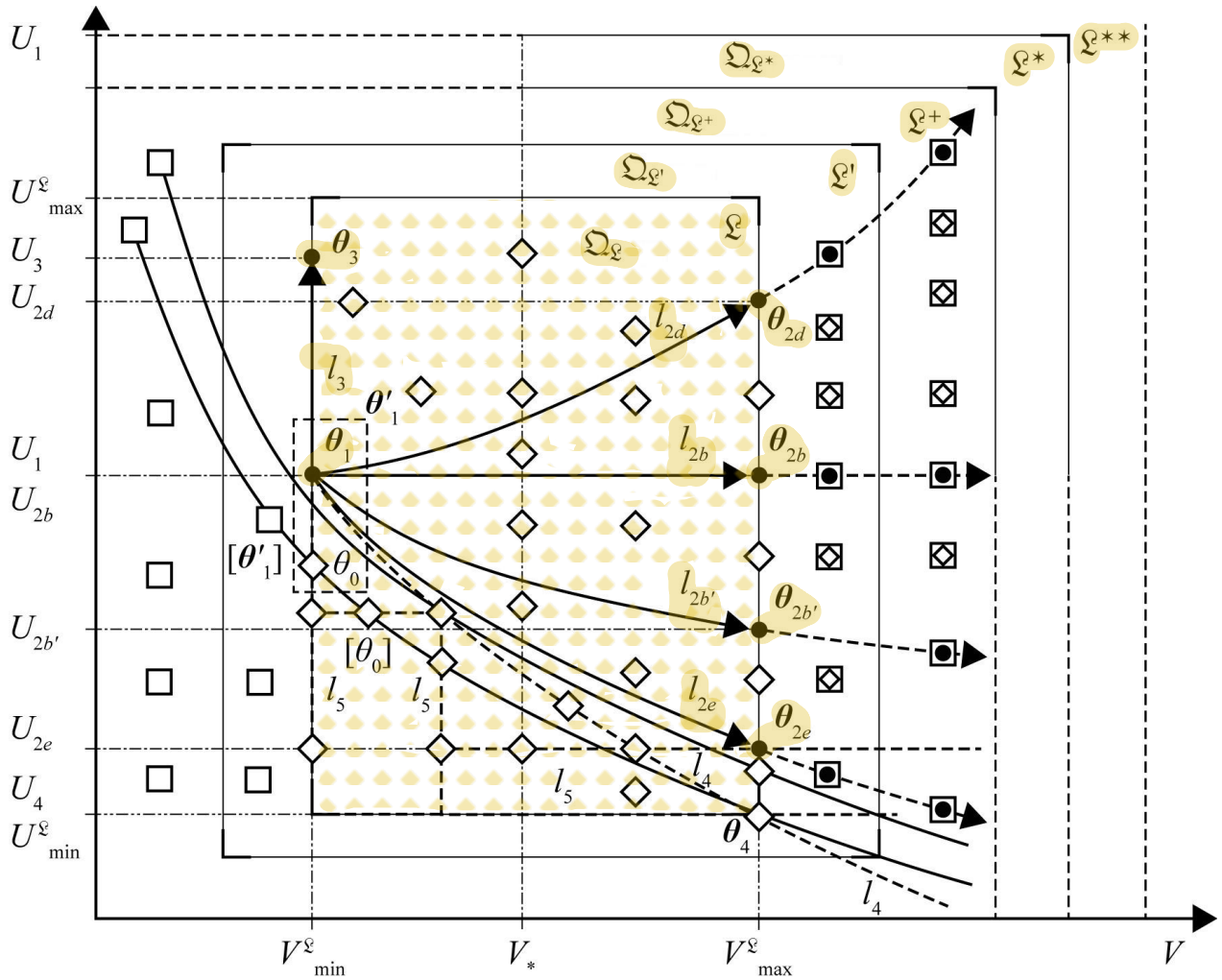
(also irreversible)

then model the consistently

inferred/inferable *PA-FORMULAS*.

2. Autoreference and Caratheodory

Any adiabatic trajectory $l_{\Omega_{\mathcal{L}}}$ is defined by its complement $\Omega_{\mathcal{L}} - l_{\Omega_{\mathcal{L}}}$, within its definition space $\Omega_{\mathcal{L}}$ /system \mathcal{L} and, the $\Omega_{\mathcal{L}} - l_{\Omega_{\mathcal{L}}}$ is not reachable within $l_{\Omega_{\mathcal{L}}}$ itself.



For the trajectory $l_{\Omega_{\mathcal{L}}}$ could 'prove' - by itself - its own adiabatic property $[d]Q_{\text{Ext}}=0$ it should have to contain its own definition as its own status!

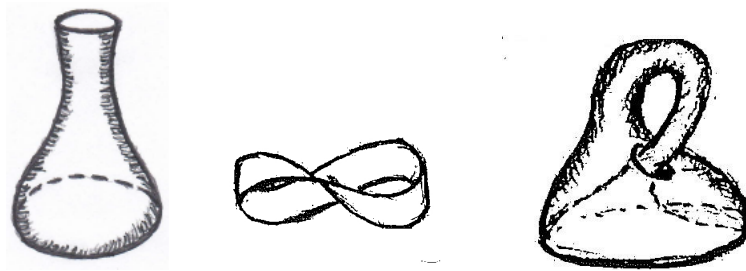
The $l_{\Omega_{\mathcal{L}}}$ would be autoreferential/autoconstructive: as the adiabatic one to construct, not adiabatically, $dQ_{\text{Ext}} \neq 0$, the adiabatic, $[d]Q_{\text{Ext}}=0$, spaces $\Omega_{\mathcal{L}}/\Omega_{\mathcal{L}} - l_{\Omega_{\mathcal{L}}}$ and define, by this way, itself - as its own status.

It is the Klein bottle building from the original's inside.

The original's outer surface defines its inner surface, which is now the model of the trajectory $l_{\Omega_{\mathcal{E}}}$ and its outer surface is the model of the $\Omega_{\mathcal{E}} - l_{\Omega_{\mathcal{E}}}$.

Within the inner surface we want to prove its internality by reaching its outer surface with an inner? curve.

Within the Klein bottle, **constructible only** by the outer manipulation with the original one, one curve is possible but, it is crossing, contradictorily against the original bottle, its inner and outer space simultaneously.⁶



Only we, as the outer constructors of the original bottle know its all properties.

The adiabatic trajectory $l_{\Omega_{\mathcal{E}}}$ does not contain itself as the object of its own

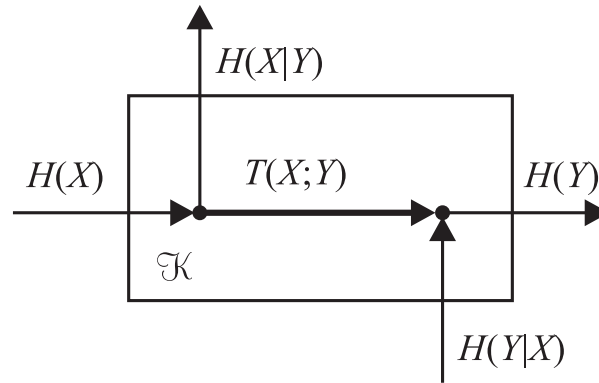
- it does not know its own properties -
and it itself, by its means, does not reveal them.

Only we, as the outer constructors of the $\Omega_{\mathcal{E}}/\mathcal{E} // \mathcal{P}$ know their **adiabaticity** // **consistency**.

⁶The original bottle should be the autoconstructive/autoreferential. The bottle should construct itself by stepping out from itself and form itself from its material as the two surfaces again and, by not with their own means.

3. Information Transfer Channel

$$\mathcal{K} \stackrel{\text{Def}}{=} [\mathbf{X}, \varepsilon, \mathbf{Y}]$$



$\mathbf{X} \stackrel{\text{Def}}{=} [\mathbf{A}, \mathbf{p}_X(\cdot)]$ - the *transmitter* of *input* messages $\mathbf{x} \in \mathbf{A}^+$

$\mathbf{Y} \stackrel{\text{Def}}{=} [\mathbf{B}, \mathbf{p}_Y(\cdot)]$ - the *receiver* of *output* messages $\mathbf{y} \in \mathbf{B}^+$,⁷

ε - the *maximal probability* of $\mathbf{y} = \mathbf{b}$ erroneous for $\mathbf{x} = \mathbf{a}$,

$\mathbf{p}_X(\cdot), \mathbf{p}_Y(\cdot)$ - the *probability distribution* on \mathbf{A} and \mathbf{B} ,
 $\mathbf{A} = \mathbf{B} = \mathbf{T}$

$\mathbf{H}(\mathbf{X}), \mathbf{H}(\mathbf{Y})$ - the *input/output information entropies*⁸

$$\mathbf{H}(\mathbf{X}) \stackrel{\text{Def}}{=} - \sum_{\mathbf{A}} \mathbf{p}_X(\cdot) \ln \mathbf{p}_X(\cdot)$$

$$\mathbf{H}(\mathbf{Y}) \stackrel{\text{Def}}{=} - \sum_{\mathbf{B}} p_Y(\cdot) \ln p_Y(\cdot)$$

$$i(\cdot) = -\ln(\cdot), \quad i(\cdot|\cdot) = -\ln(\cdot|\cdot)$$

⁷ \mathbf{A}, \mathbf{B} - a finite *alphabets* of elements \mathbf{x} of \mathbf{X} and \mathbf{y} of \mathbf{Y} .

⁸*Shannon entropies* - average amounts of information in any $\mathbf{x} \in \mathbf{A}$ and $\mathbf{y} \in \mathbf{B}$.

$\mathbf{H}(\mathbf{X}|\mathbf{Y})$, $H(Y|X)$ - the *loss/noise entropy*

$$\mathbf{H}(\mathbf{X}|\mathbf{Y}) \stackrel{\text{Def}}{=} - \sum_A \sum_B \mathbf{p}_{\mathbf{X},\mathbf{Y}}(\cdot, \cdot) \ln \mathbf{p}_{\mathbf{X}|\mathbf{Y}}(\cdot|\cdot)$$

$$H(Y|X) \stackrel{\text{Def}}{=} - \sum_A \sum_B p_{X,Y}(\cdot, \cdot) \ln p_{Y|X}(\cdot|\cdot)$$

For the *transinformation* $\mathbf{T}(\mathbf{X}; \mathbf{Y})$, $T(Y; X)$ ⁹

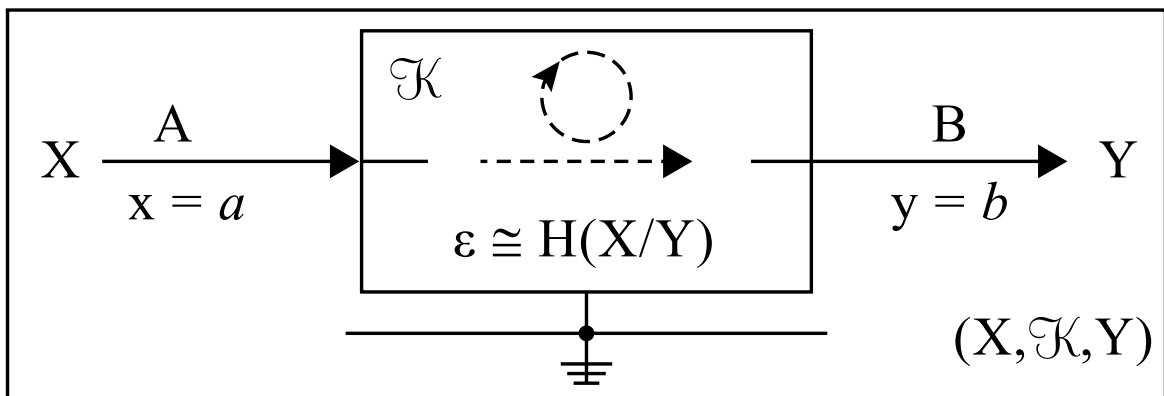
$$\mathbf{T}(\mathbf{X}; \mathbf{Y}) \stackrel{\text{Def}}{=} \mathbf{H}(\mathbf{X}) - \mathbf{H}(\mathbf{X}|\mathbf{Y})$$

$$T(Y; X) \stackrel{\text{Def}}{=} H(Y) - H(Y|X)$$

the **channel equation** is valid

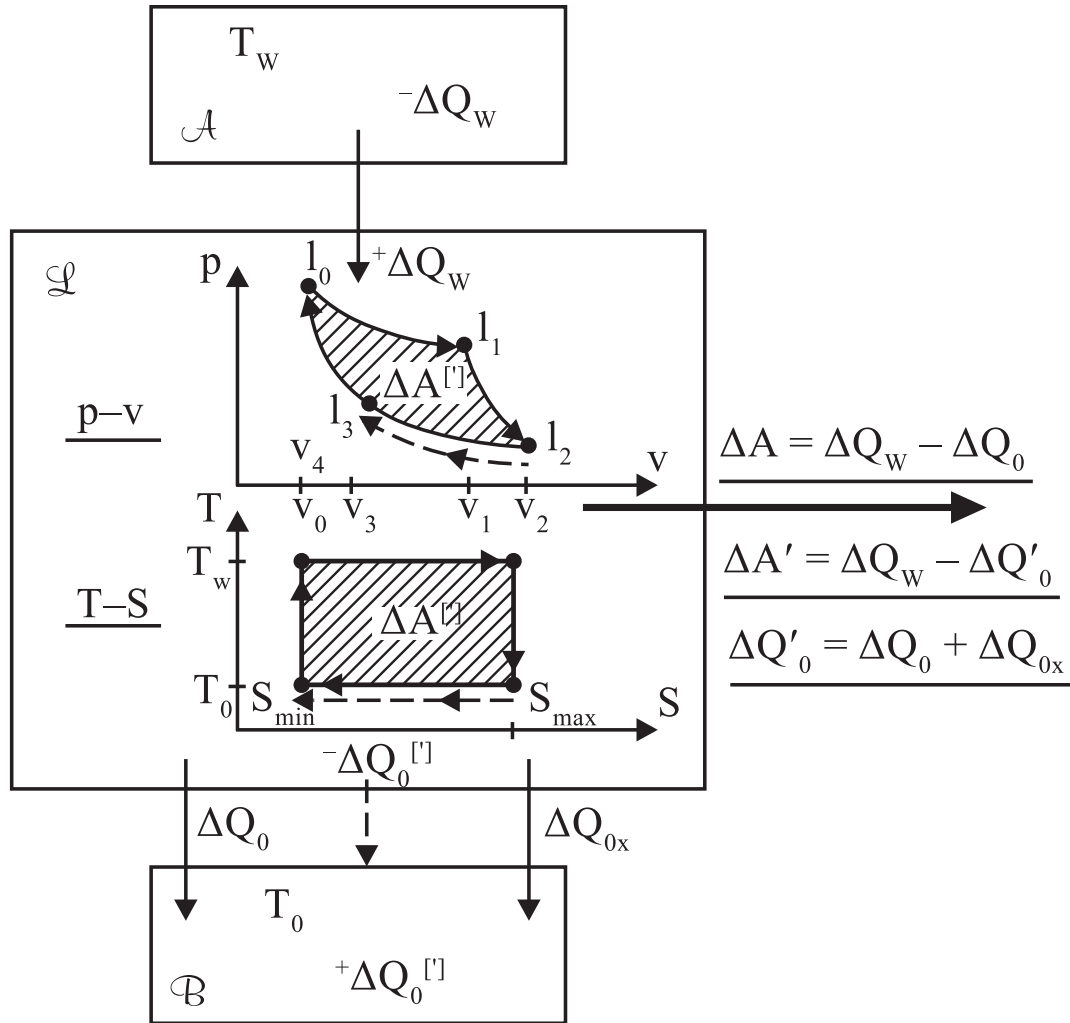
$$\mathbf{H}(\mathbf{X}) - \mathbf{H}(\mathbf{X}|\mathbf{Y}) = H(Y) - H(Y|X)$$

$(\mathbf{X}, \mathcal{K}, \mathbf{Y})$ - *Shannon Transfer Chain*



⁹Also it is valid for the information \mathbf{i} with the probability $\mathbf{p}(\cdot)$, $\mathbf{i}_X + \mathbf{i}_{Y|X} = \mathbf{i}_Y + \mathbf{i}_{X|Y}$, $\mathbf{i} = -\ln \mathbf{p}(\cdot)$.

4. Carnot Machine



l_0 - l_1 : *isothermal exp.* transfers the heat ΔQ_W from \mathcal{A} to \mathcal{L} ,
the work $\Delta A_{0,1} = \Delta Q_W$ is **given** at T_W

l_1 - l_2 : *adiabatic exp.* cools \mathcal{L} from T_W to T_0 , the work
 $\Delta A_{1,2} = -\Delta U$ is **given** from the internal energy U of \mathcal{L}

l_2 - l_3 : *isothermal comp.* transfers the heat $\Delta Q_0 < \Delta Q_W$ ¹⁰
from \mathcal{L} to \mathcal{B} at T_0 , **consumes** the work $-\Delta A_{2,3} < \Delta A_{0,1}$

l_3 - l_0 : *adiabatic comp.* heats \mathcal{L} from T_0 to T_W , $\Delta U > 0$,
and **consumes** the work $-\Delta A_{3,4} = \Delta A_{1,2}$

¹⁰In the reversible Carnot Cycle is $\Delta Q_{0x} = 0$, no *production* of (positive) *noise* heat, $\Delta Q_{0x} > 0$, arises.

The resulting output work for a reversible Carnot Cycle \mathcal{O} is

$$\Delta A = \Delta Q_w - |\Delta Q_0|$$

$$\Delta A = \Delta A_{l_0-l_1} + \Delta A_{l_1-l_2} + \Delta A_{l_2-l_3} + \Delta A_{l_3-l_0}$$

Kelvin's form of the II. P. T. for a reversible case is

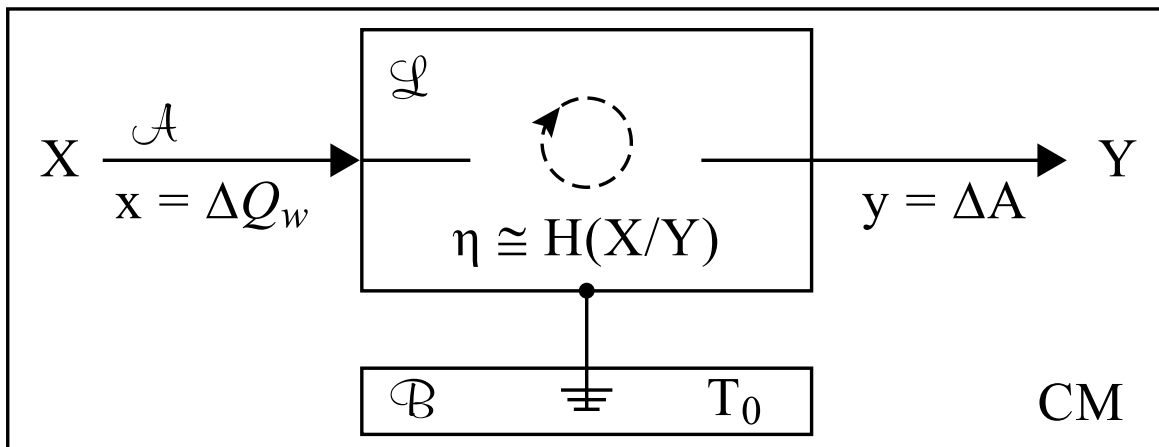
$$\sum_{i \in [W, 0]} \frac{\Delta Q_i}{T_i} \triangleq \oint_{\mathcal{O}} \frac{\delta Q(\Theta)}{\Theta} = 0$$

Thomson-Planck's formulation of the II. P. T. says:

It is impossible to construct a heat cycle transforming all heat delivered to the medium \mathcal{L} (going through this cycle) into the equivalent amount of the mechanical work ΔA .

The I. P.T. is valid \Leftrightarrow the transformation efficiency η_{\max} :

$$\eta_{\max} \stackrel{\text{Def}}{=} \frac{\Delta A}{\Delta Q_w} = \frac{\Delta Q_w - |\Delta Q_0|}{\Delta Q_w} = \frac{T_w - T_0}{T_w} < 1$$



5. Carnot Cycle and Noiseless Information Transfer

Recording/transmitting/computing an information ΔI at the temperature Θ requires the energy ΔW

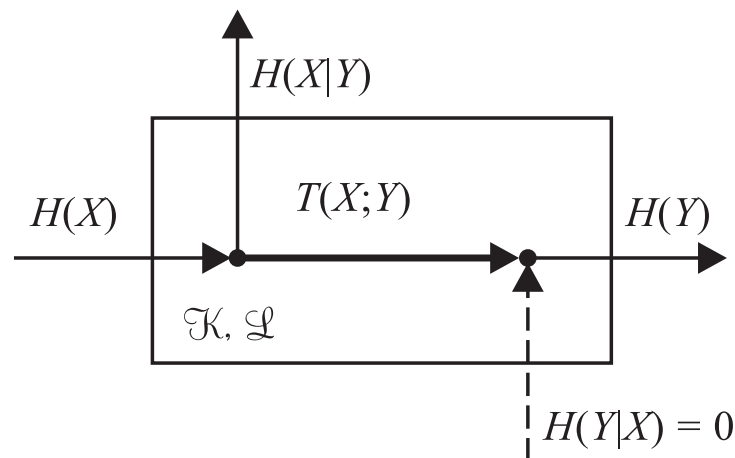
$$\Delta W \geq k \cdot \Theta \cdot \Delta I, \quad \text{now} \quad \Delta W \triangleq \Delta Q_W$$

The changes of the entropies of the medium \mathcal{L} with \mathcal{O} are now considered informationally on a \mathcal{K}

$$\begin{aligned} H(\mathbf{X}) &\stackrel{\text{Def}}{=} \frac{\Delta Q_W}{kT_W}, & H(\mathbf{Y}|\mathbf{X}) &\stackrel{\text{Def}}{=} \mathbf{0} \\ H(\mathbf{Y}) &\stackrel{\text{Def}}{=} \frac{\Delta A}{kT_W} = \frac{\Delta Q_W - \Delta Q_0}{kT_W} \\ &= \frac{\Delta Q_W}{kT_W} \cdot \eta_{\max} = H(\mathbf{X}) \cdot \eta_{\max} \triangleq \Delta I \end{aligned}$$

11

$$H(\mathbf{Y}) - H(\mathbf{Y}|\mathbf{X}) = H(\mathbf{X}) - H(\mathbf{X}|\mathbf{Y})$$

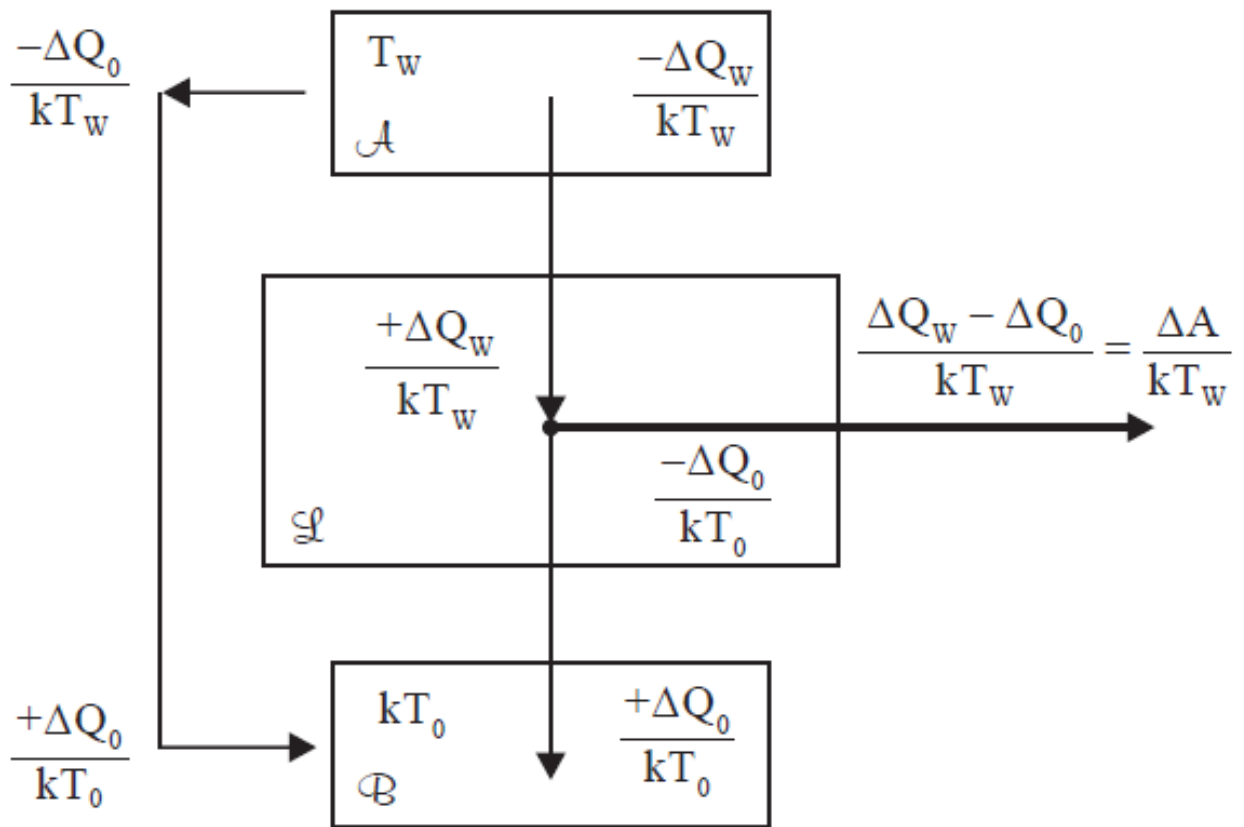


With our thermodynamic substitutions we gain:

¹¹ $H(X) \geq H(Y) = T(X;Y) = \Delta I \geq 0$; the *information* form of the II. P.T. is implied for the reversible case; Brillouin, Landauer, Gershenfeld, Bennet.

$$\frac{\Delta Q_W}{kT_W} \cdot \eta_{\max} - 0 = \frac{\Delta Q_W}{kT_W} - \mathbf{H}(\mathbf{X}|\mathbf{Y})$$

$$\mathbf{H}(\mathbf{X}|\mathbf{Y}) = \frac{\Delta Q_W}{kT_W} \cdot (1 - \eta_{\max}) = \frac{\Delta Q_0}{kT_W}$$



The change $\Delta S_{\mathcal{AB}}$ within the change $\Delta S_{\mathcal{C}}$
of the global CM 's heat entropy $S_{\mathcal{C}}$
within its subsystem \mathcal{AB} is

$$\Delta S_{\mathcal{AB}} = -\frac{\Delta Q_0}{T_W} + \frac{\Delta Q_0}{T_0} = \frac{\Delta Q_0}{T_0} \cdot \eta_{\max} = \frac{\Delta Q_W}{T_W} \cdot \eta_{\max}$$

The change $\Delta S_{\mathcal{L}}$ of the heat entropy $S_{\mathcal{L}}$ within the change $\Delta S_{\mathcal{C}}$ of the whole heat entropy $S_{\mathcal{C}}$ of the CM ¹² is

$$\Delta S_{\mathcal{L}} = \oint_{\mathcal{O}} \frac{\delta Q}{T} = \frac{\Delta Q_{\mathcal{W}}}{T_{\mathcal{W}}} - \frac{\Delta Q_0}{T_0} = 0$$

The resultant change $\Delta S_{\mathcal{C}}$ of CM and the output ΔI is

$$\Delta S_{\mathcal{C}} = \Delta S_{\mathcal{L}} + \Delta S_{AB} = \frac{\Delta Q_{\mathcal{W}}}{T_{\mathcal{W}}} \cdot \eta_{\max} = k \cdot \Delta I = k \cdot H(Y)$$

The Brillouin's *extended* form of the *II. P.T.* is valid¹³

$$d(S_{\mathcal{C}} - k \cdot I) \geq 0$$

$$\begin{aligned} \Delta S_{\mathcal{C}} - k \cdot T(X; Y) &= k \cdot H(X) \cdot (\eta_{\max} - \eta_{\max}) \\ \Delta S_{\mathcal{C}} - k \cdot \Delta I &= 0 \quad \Delta(S_{\mathcal{C}} - k \cdot I) = 0 \end{aligned}$$

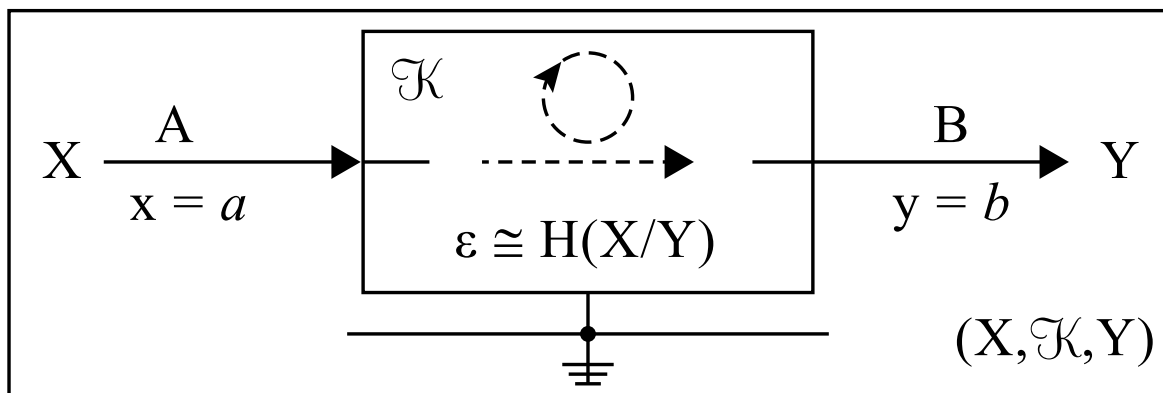
¹²Or of the whole system in which the CM is running.

¹³The *information* member I does not exist in the traditional (differential) formulation of this theorem; it is $dS \geq 0$ only.

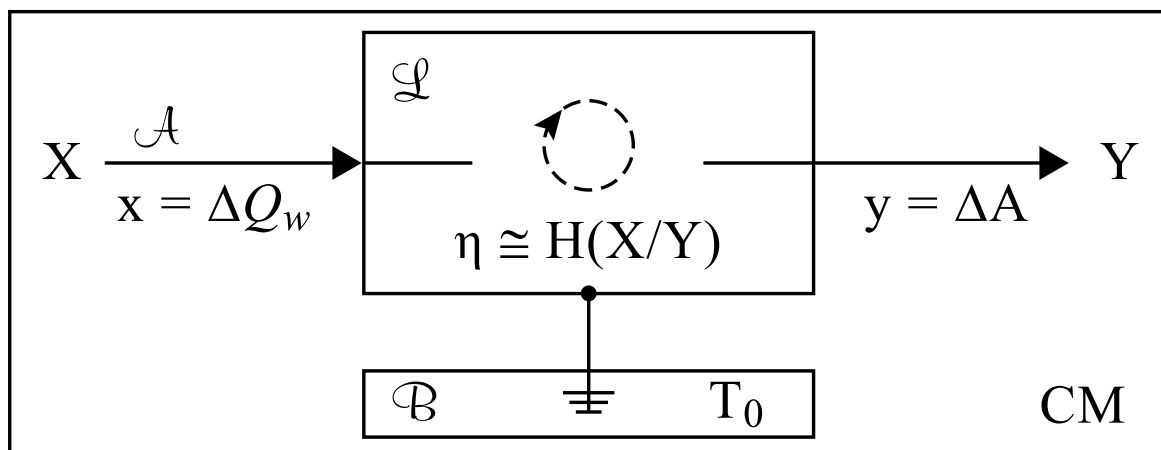
- ★ The *reversible* Carnot Cycle \mathcal{O}
- the **medium** \mathcal{L} going through the \mathcal{O}
- the **whole** CM

work as *thermodynamic* models of

- ★ the **information transfer process** \mathcal{T} *without noise*,¹⁴
- the **channel** \mathcal{K} with its transfer process \mathcal{T}
- the *Shannon Transfer Chain* $(\mathbf{X}, \mathcal{K}, \mathbf{Y})$



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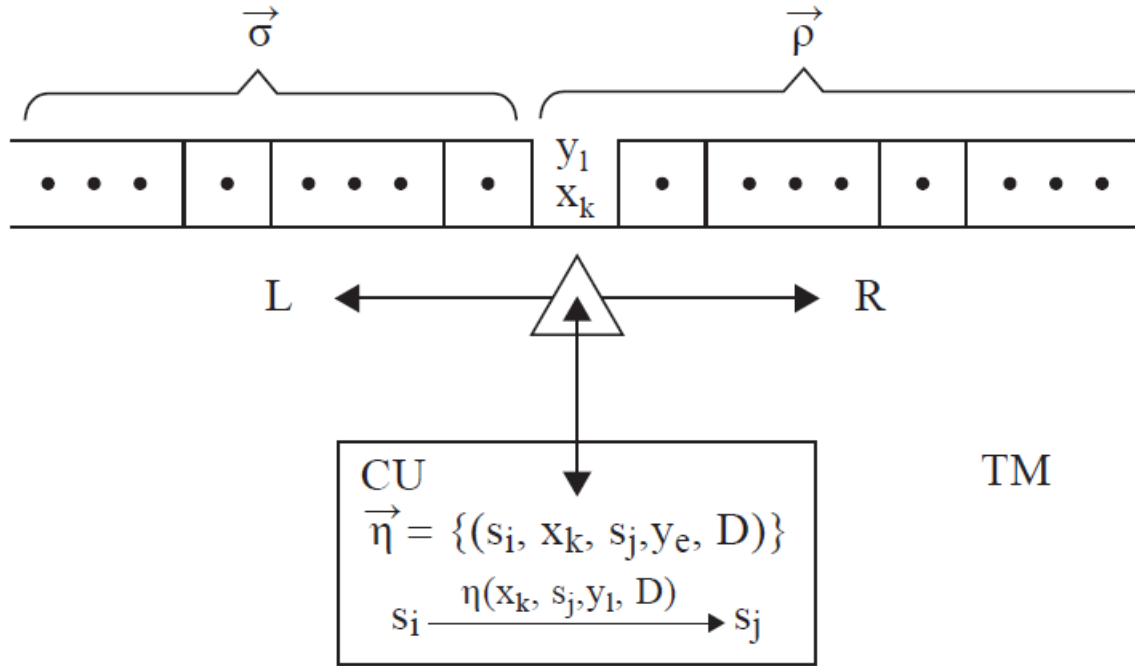
¹⁴ $H(\mathbf{Y}|\mathbf{X}) = 0$

6. Turing Machine

Turing Machine TM - driven by the program $\vec{\eta}$

$$\vec{\eta} = (\eta_p)_{p=1}^{p \in \mathbb{N}} = [(\mathbf{s}_i, \mathbf{x}_k, \mathbf{s}_j, \mathbf{y}_1, \mathbf{D})_p]_{p=1}^{p \in \mathbb{N}}, \quad |\vec{\eta}| \in \mathbb{N}$$

$$\eta_{[\cdot]} = (\mathbf{s}_{i[\cdot]}, \mathbf{x}_{k[\cdot]}, \mathbf{s}_{j[\cdot+1]}, \mathbf{y}_{1[\cdot]}, \mathbf{D})$$



\mathbf{s}_i - the *status* of the \mathbf{CU}_{TM} in the actual step $\mathbf{p} \in \mathbb{N}$

\mathbf{x}_k - the *input symbol* on the *input-output tape* in the step \mathbf{p}

\mathbf{y}_1 - the *output symbol overwriting \mathbf{x}_k* in the step \mathbf{p}

\mathbf{s}_j - the *defined \mathbf{CU}_{TM} 's status* for the step $\mathbf{p} + 1$

\mathbf{D} - the \mathbf{CU}_{TM} *read-write head moving Left/Right* after \mathbf{y}_1 has been written instead of \mathbf{x}_k in the step \mathbf{p}

$$[\mathbf{y}_1, \mathbf{x}_k \in \mathbf{T} = \mathbf{A} = \mathbf{B}]$$

$(\vec{\sigma}, \mathbf{s}_i, \vec{\varrho}) / (\mathbf{x}_k, \mathbf{s}_i, \mathbf{y}_1)$ - the *TM's/ \mathbf{CU}_{TM} 's configurations*

With the I/O transformations $\mathbf{x}_k \longrightarrow \mathbf{y}_l$,
the CU_{TM} 's states' *transitions* are performed,

$$\mathbf{S}_{i_p} \xrightarrow{(\mathbf{x}_{k_p}, \mathbf{y}_{l_p}, \mathbf{D}_p)} \mathbf{S}_{j_{p+1}}$$

defining the regular grammar and the language $\text{L}_{\text{CU}_{TM}}$

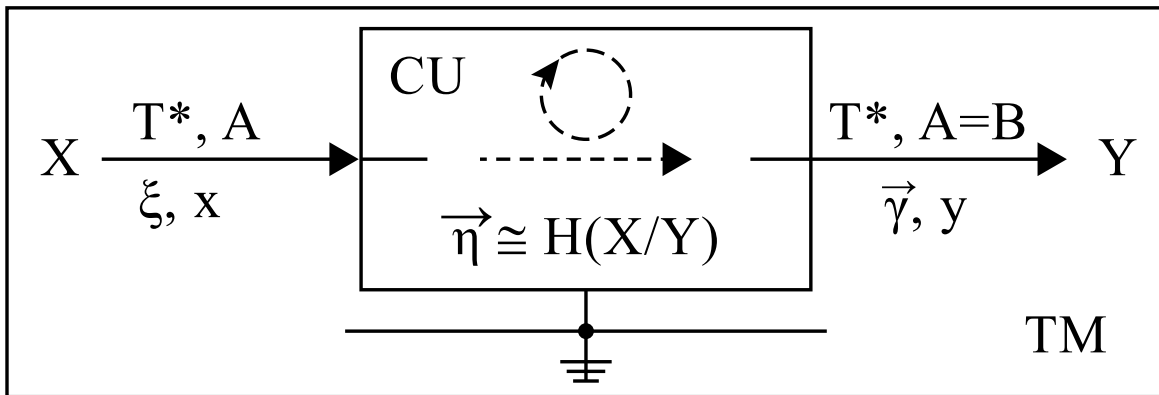
$$\mathbf{S}_{i_p} \longrightarrow (\mathbf{x}_{k_p}, \mathbf{y}_{l_p}, \mathbf{D}_p) \mathbf{S}_{j_{p+1}}$$

$$\text{L}_{\text{CU}_{TM}} = \{(\mathbf{x}_{k_p}, \mathbf{y}_{l_p}, \mathbf{D}_p)\}_{p=1}^{p=\text{last}}$$

and the regular language L_{TM} of the configurations
which the TM has gone through so far¹⁵

$$\mathbf{S}_{i_p} \longrightarrow (\vec{\sigma}_p, \mathbf{S}_{i_p}, \vec{\rho}_p) \mathbf{S}_{j_{p+1}}$$

$$\text{L}_{TM} = \{(\vec{\sigma}_p, \mathbf{S}_{i_p}, \vec{\rho}_p)\}_{p=1}^{p=p_{\text{last}}}$$



¹⁵Terminal symbols $\mathbf{T} = \{I, B\}$ • the instruction (s_i, x_k, s_j, y_l, D) • the configuration $(\vec{\sigma}, s_{[i]}, \vec{\rho})$ • the configuration type $(\varepsilon [\sigma, s_{[i]}, \rho] \varepsilon)$ • $\mathbf{X} = (\vec{\sigma}, s_{[i]}, \vec{\rho}) \triangleq (\mathbf{B}\sigma, s_{[i]}, \rho\mathbf{B})$ the general configuration type, e.g. $\mathbf{B}\bar{1}\bar{B}s_{[i]}\bar{1}\bar{B}$. Also $\mathbf{S}_{i_p} \longrightarrow (\mathbf{S}_{i_p}, \mathbf{x}_{k_p}, \mathbf{S}_{j_{p+1}}, \mathbf{y}_{l_p}, \mathbf{D}_p) \mathbf{S}_{j_{p+1}}$, $\text{L}'_{TM} = \{(\mathbf{S}_{i_p}, \mathbf{x}_{k_p}, \mathbf{S}_{j_{p+1}}, \mathbf{y}_{l_p}, \mathbf{D}_{[.]})\}_{p=1}^{p=p_{\text{last}}}$, instructions have been performed yet, [14, 19, 18].

7. Inference as Information Transfer

The right *TM*'s program $\vec{\eta}$ ¹⁶ generates the resultative *TM*'s configuration sequence and

similar is valid for the information transfer acts and the $(\mathbf{X}, \mathcal{K}, \mathbf{Y})$'s configurations.

Now, the inferred *CLAIM* \mathbf{a}_i is the last member of the input *FORMULAE* chain $\vec{\mathbf{x}}$ ¹⁷ and the \mathbf{a}_i 's $\mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ inference by *Modus Ponens* is realized as the \mathbf{y} 's information transfer \mathcal{T} in \mathcal{K} .

$$[\vec{\mathbf{x}}|\vec{\mathbf{y}}] = [\mathbf{a}_0, \mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_{i-1}] \sqsubset [\vec{\mathbf{x}}]$$

$$[\vec{\mathbf{x}}] = [\mathbf{a}_0, \mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_{i-1}; \mathbf{a}_i]$$

$$[\mathbf{y}] = \mathbf{a}_i \sqsubset [\vec{\mathbf{x}}]$$

Entropies for this \mathbf{a}_i 's $\mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ inference from $\vec{\mathbf{x}}$ realized by the $\mathbf{y} = \mathbf{a}_i$ transfer from $\vec{\mathbf{x}}$ through a \mathcal{K} in its status $[\vec{\mathbf{x}}|\vec{\mathbf{y}}]$ are:

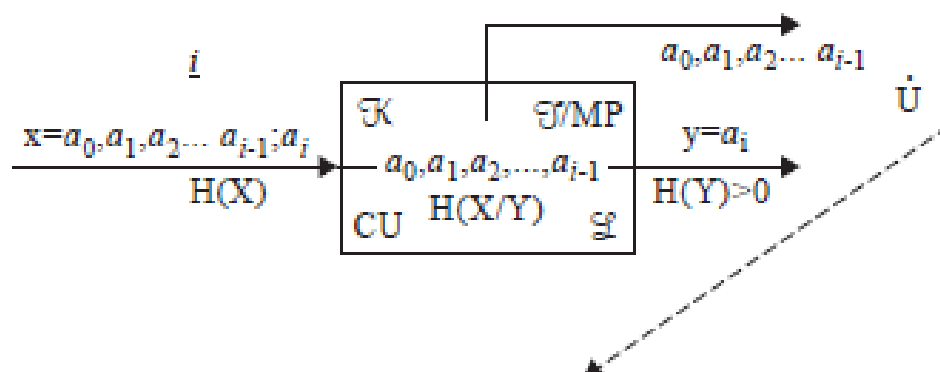
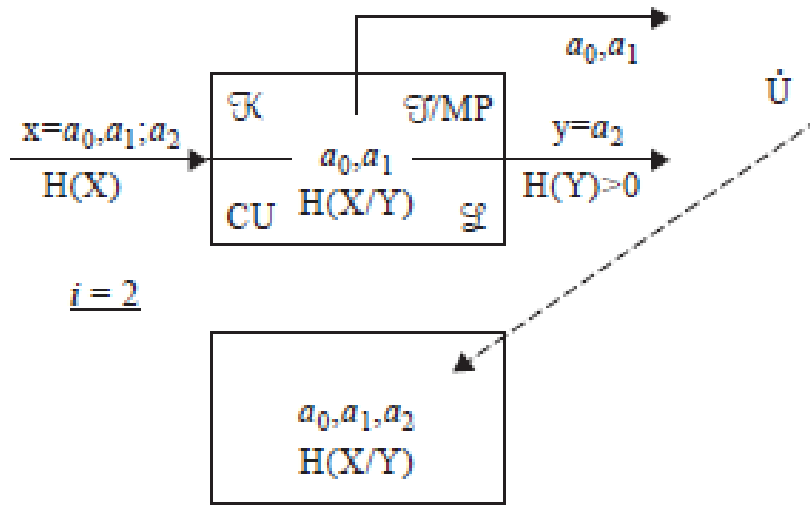
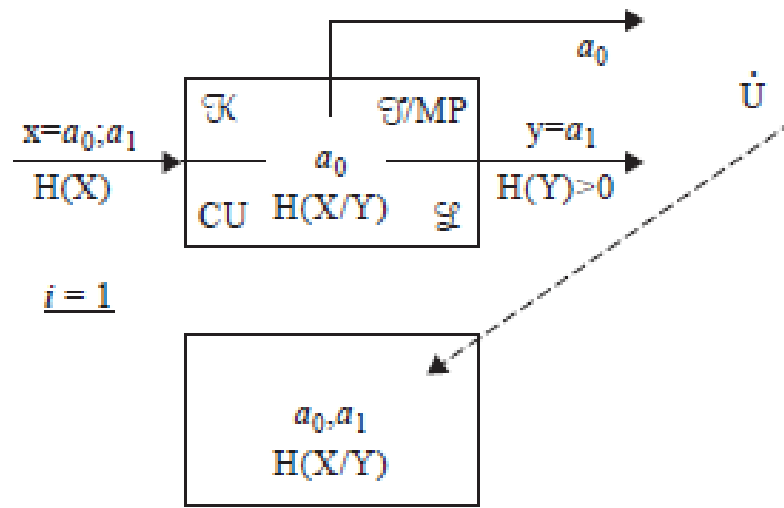
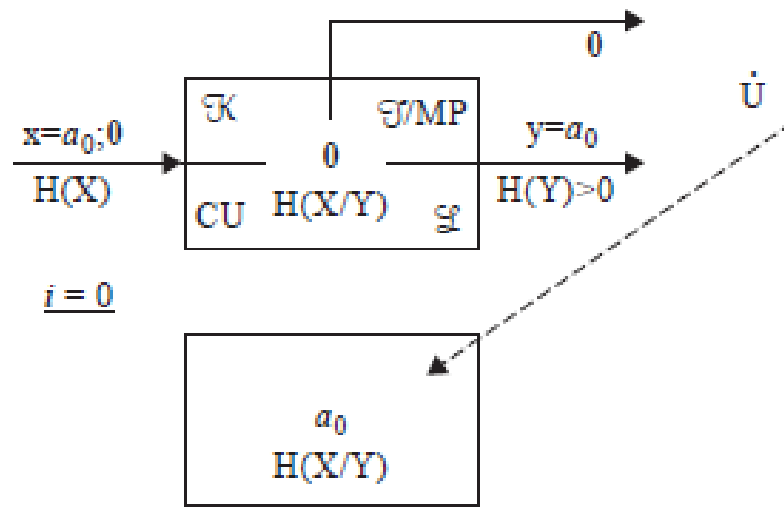
$$\mathbf{H}(\mathbf{X}|\mathbf{Y}) \cong \mathbf{H}(\mathbf{a}_0, \mathbf{a}_1, \dots, \mathbf{a}_{i-1}) = \mathbf{H}(\vec{\mathbf{x}}|\vec{\mathbf{y}}), \quad \mathbf{H}(\mathbf{Y}) = \mathbf{H}(\mathbf{y})$$

$$\mathbf{H}(\mathbf{X}) \cong \mathbf{H}(\mathbf{a}_0, \mathbf{a}_1, \dots, \mathbf{a}_{i-1}; \mathbf{a}_i) = \mathbf{H}(\vec{\mathbf{x}})$$

$$\mathbf{H}(\mathbf{Y}) \cong \mathbf{H}(\mathbf{a}_0, \mathbf{a}_1, \dots, \mathbf{a}_{i-1}; \mathbf{a}_i) - \mathbf{H}(\mathbf{a}_0, \mathbf{a}_1, \dots, \mathbf{a}_{i-1})$$

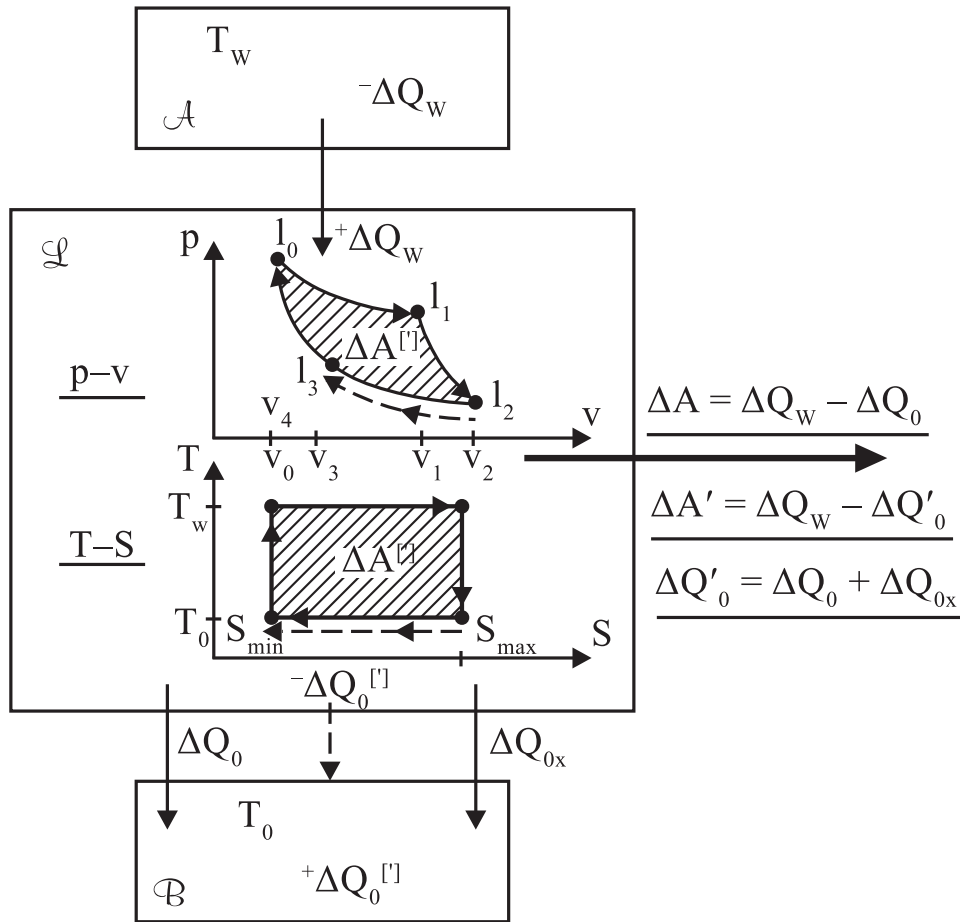
¹⁶In our thermodynamic analogy following the Caratheodory II. P.T.

¹⁷With the Gödel number \mathbf{x} ,



8. Carnot Cycle, Automata and Information Transfer

Any CM , now with \mathcal{O} , is describable as an **automaton**¹⁸



with the resulting regular grammar and language L ,

$$\{\theta_{l_0}^i \rightarrow (\Delta A^i) \theta_{l_0}^i\}, \quad L = \{\Delta A^i\} \quad \text{and, as}$$

for \mathcal{K}

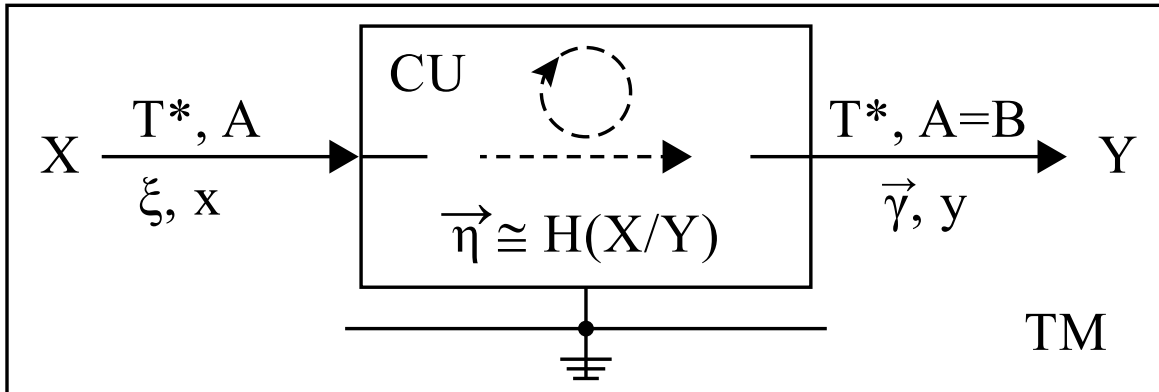
$$H^i(\mathbf{X}) = \frac{\Delta A_{l_0-l_1}^i}{kT_w^i}, \quad H^i(\mathbf{X}|\mathbf{Y}) = \frac{|\Delta A_{l_2-l_3}^i|}{kT_w^i}$$

$$H^i(\mathbf{Y}) = H^i(\mathbf{X}) - H^i(\mathbf{X}|\mathbf{Y}) = \frac{\Delta A^i}{kT_w^i} > 0$$

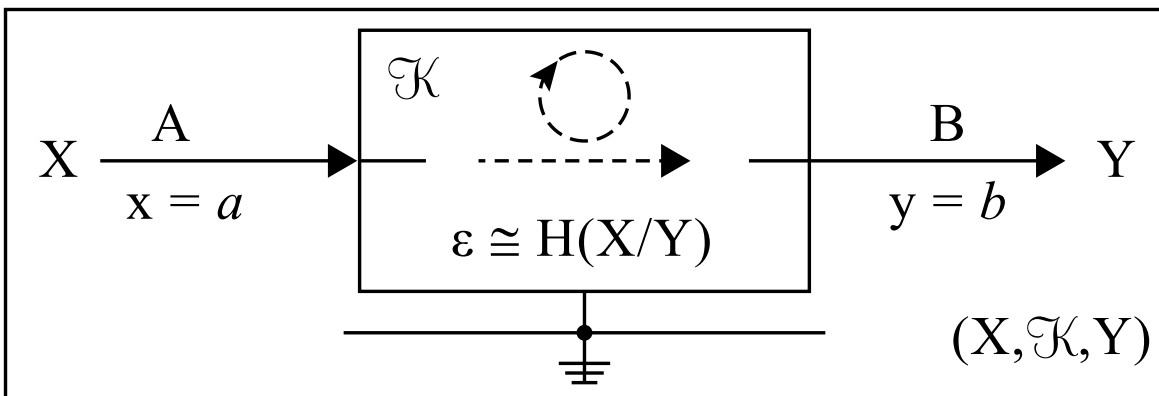
$$\Delta A^i = \Delta A_{l_0-l_1}^i + \Delta A_{l_1-l_2}^i + \Delta A_{l_2-l_3}^i + \Delta A_{l_3-l_0}^i$$

¹⁸Now the Moore's but not only. Information transmission (not cyclical or cyclical) or the heat energy transformation (not cyclical or cyclical) is also describable by the terminology of *regular grammars* and *finite automata*.

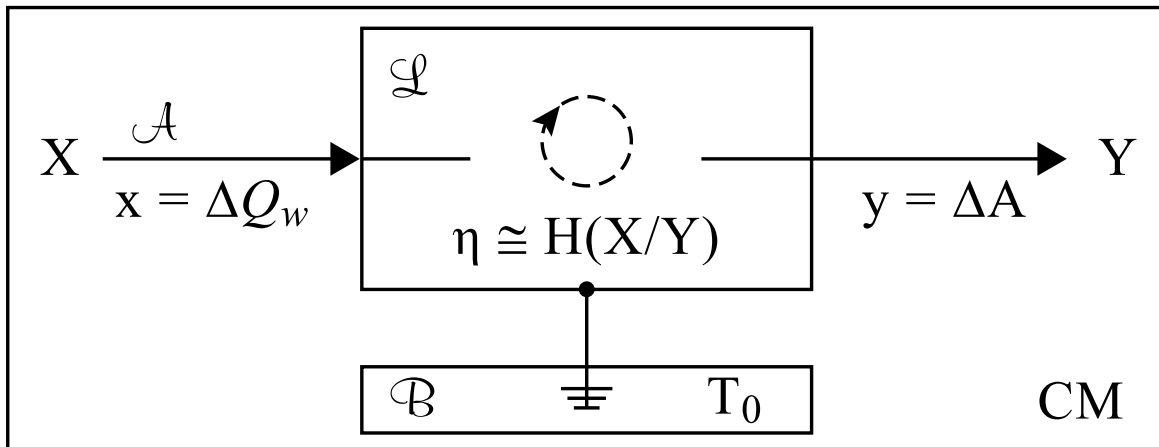
CU_{TM} 's step $\mathbf{p} \cong \mathcal{K}$'s transfer act $\mathbf{i} \cong \mathcal{L}$'s cycle \mathcal{O} run \mathbf{i}
 $\mathcal{T}_{\mathcal{P}_A}$ Inference \cong Message Transfer \cong Heat Transformation
 states: $\text{CU}_{TM} \cong \mathcal{K} \cong \mathcal{L}$
 config: $(\vec{\sigma}, \mathbf{s}_i, \vec{p}) \cong [(\mathbf{X})^i, \mathbf{X}^i | \mathbf{Y}^i, (\mathbf{X})^{i+1}] \cong [\sum Q_w^i, (\mathbf{p}^i, \mathbf{V}^i, \mathbf{T}^i)_{\mathcal{L}}, Q_w - \sum Q_w^i]$
 $TM \cong (\mathbf{X}, \mathcal{K}, \mathbf{Y}) \cong CM$



\cong



\cong



The TM 's, $(\mathbf{X}, \mathcal{K}, \mathbf{Y})$'s, CM 's runs are considered in isolated systems for X , Y and $X|Y$ energies.

9. Resultativity, Adiabaticity, Consistency

The states' $\theta_{[\cdot]}^{\mathcal{L}}$ changes in the adiabatic system $\mathcal{L}/\Omega_{\mathcal{L}}$,
along the trajectories $l_{\Omega_{\mathcal{L}}}$ are expressible **regularly**:

l_{2b} <i>isothermic irreversible,</i>	$\theta_1^{\mathcal{L}} \rightarrow \Delta A_{1,2e} \theta_{2e}^{\mathcal{L}}$
$l_{2b'}$ <i>adiabatic irreversible,</i>	$\theta_1^{\mathcal{L}} \rightarrow \Delta A_{1,3} \theta_3^{\mathcal{L}}$
l_{2d} <i>izobaric irreversible,</i>	$\theta_1^{\mathcal{L}} \rightarrow \Delta A_{1,2b} \theta_{2b}^{\mathcal{L}}$
l_{2e} <i>izentropic,</i>	$\theta_1^{\mathcal{L}} \rightarrow \Delta A_{1,2b'} \theta_{2b'}^{\mathcal{L}}$
l_3 <i>izochoric irreversible,</i>	$\theta_1^{\mathcal{L}} \rightarrow \Delta A_{1,2d} \theta_{2d}^{\mathcal{L}}$
	$\theta_1^{\mathcal{L}} \rightarrow \lambda \theta_1^{\mathcal{L}}$

The thermodynamic model

for **the consistent $\mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ inference**

- from its axioms/formulas having been inferred so far -

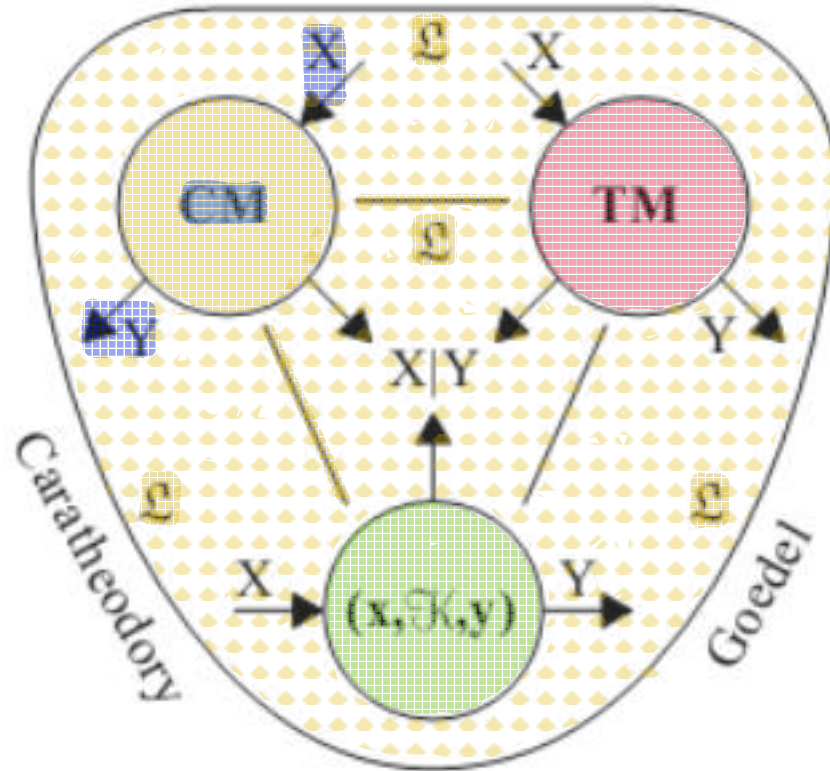
is created by the **CM's activity**, which

is modeling both the **TM** and the **(X, K, Y)**
 and which runs in the adiabatic system $\mathcal{L}/\Omega_{\mathcal{L}}$.

The **TM's**, **(X, K, Y)'s**, configurations

are then modeled by the states $\theta_i^{\mathcal{L}} \in \Omega_{\mathcal{L}}$
 of the adiabatic $\mathcal{L}/\Omega_{\mathcal{L}}$ with this modeling **CM** inside
the configuration of which,
 in fact, are creating **these states**.

II. P. T.



The \mathcal{L} 's initial imbalance starts the $\theta_{[\cdot]}^{\mathcal{L}}$ states' sequence on a trajectory $l_{\Omega_{\mathcal{L}}}$ (irreversible) and is given by the modeled

temperature difference $T_W - T_0 > 0$ on CM ,
 existence of the input message on \mathcal{K} ,
 input chain's existence on the TM 's input-output tape

\Rightarrow

These adiabatic trajectories $l_{\Omega_{\mathcal{L}}}$

now represent the norm

of the **consistency** (and resultativity)

of the **$\mathcal{P}/\mathcal{T}_{\mathcal{P}A}$ -inference/computing process**

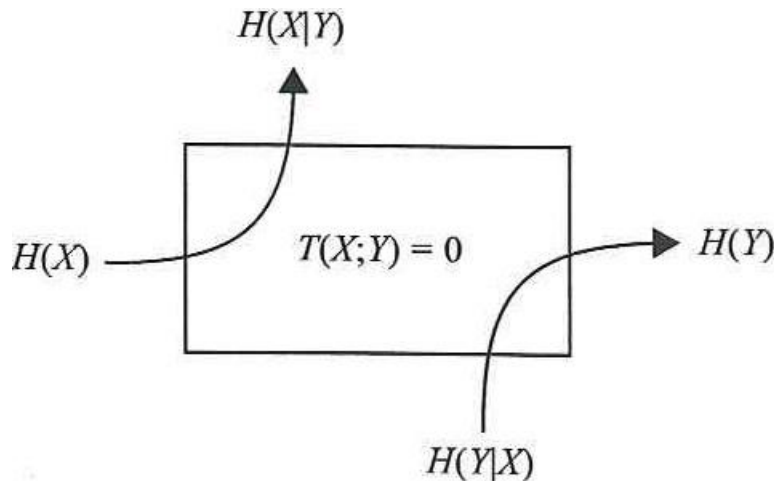
expressible also in terms

of the **information transfer/heat energy transformation**.

10. Autoreference, Information Transfer, Thermodynamics

On the *interrupted* channel \mathcal{K} is valid

$$\begin{aligned} \mathbf{H(X)} &= \mathbf{H(X|Y)} \\ \mathbf{T(X; Y)} &= \mathbf{H(X) - H(X|Y) = 0} \end{aligned}$$



The input $\mathbf{H(X)}$ is now the measure of the \mathcal{K} 's internal state - $\mathbf{H(X|Y)}$ - the output $\mathbf{H(Y)}$ is without any relation to $\mathbf{H(X)}$.

With insisting (?!?) on the information transfer through this interrupted channel \mathcal{K} , then, contradictorily, we want build the $(\mathbf{X}, \mathcal{K}, \mathbf{Y})$ with a reversible direct Carnot Machine CM , where

$$\Delta Q_W = \Delta Q_0 \triangleq \Delta Q \quad \& \quad T_W > T_0, \quad \eta_{\max} = \frac{T_W - T_0}{T_W} > 0$$

In fact we 'measure' ΔQ against ΔQ , $T_W = T_0$,¹⁹

$$\mathbf{H(X)} = \mathbf{H(X|Y)} = \frac{\Delta Q}{kT_W}, \quad \underline{\mathbf{H(Y) = 0}} \quad [= H(Y|X)]$$

¹⁹The measuring with the zero 'distance' between the measuring and measured - the Gibbs Paradox.

Our 'wish' to have the information transfer $H(Y) > 0$
 through such interrupted channel \mathcal{K} formulates
 the contradiction/paradox²⁰

- the *II. P.T.'s violation* -

$$\Delta S_{\mathcal{L}} = \oint_0 \frac{\delta Q}{T} = \frac{\Delta Q}{T_W} - \frac{\Delta Q}{T_0} = -\frac{\Delta Q}{T_0} \cdot \eta_{\max} < \underline{0} (!)$$

$$\Delta S_{AB} = \frac{\Delta Q}{T_W} \cdot \eta_{\max}$$

$$\Delta S_{\mathcal{C}} = \Delta S_{\mathcal{L}} + \Delta S_{AB} < \underline{0} (!) \quad [\text{in fact } 0 \text{ is everywhere}]$$

Our 'measuring' is now with the 0 'distance' between
 the measuring object \mathcal{K}/\mathcal{L}

and the measured object \mathbf{X}/\mathcal{A} and

we see that the equality $H(\mathbf{X}) = H(\mathbf{X}|\mathbf{Y})$ says that

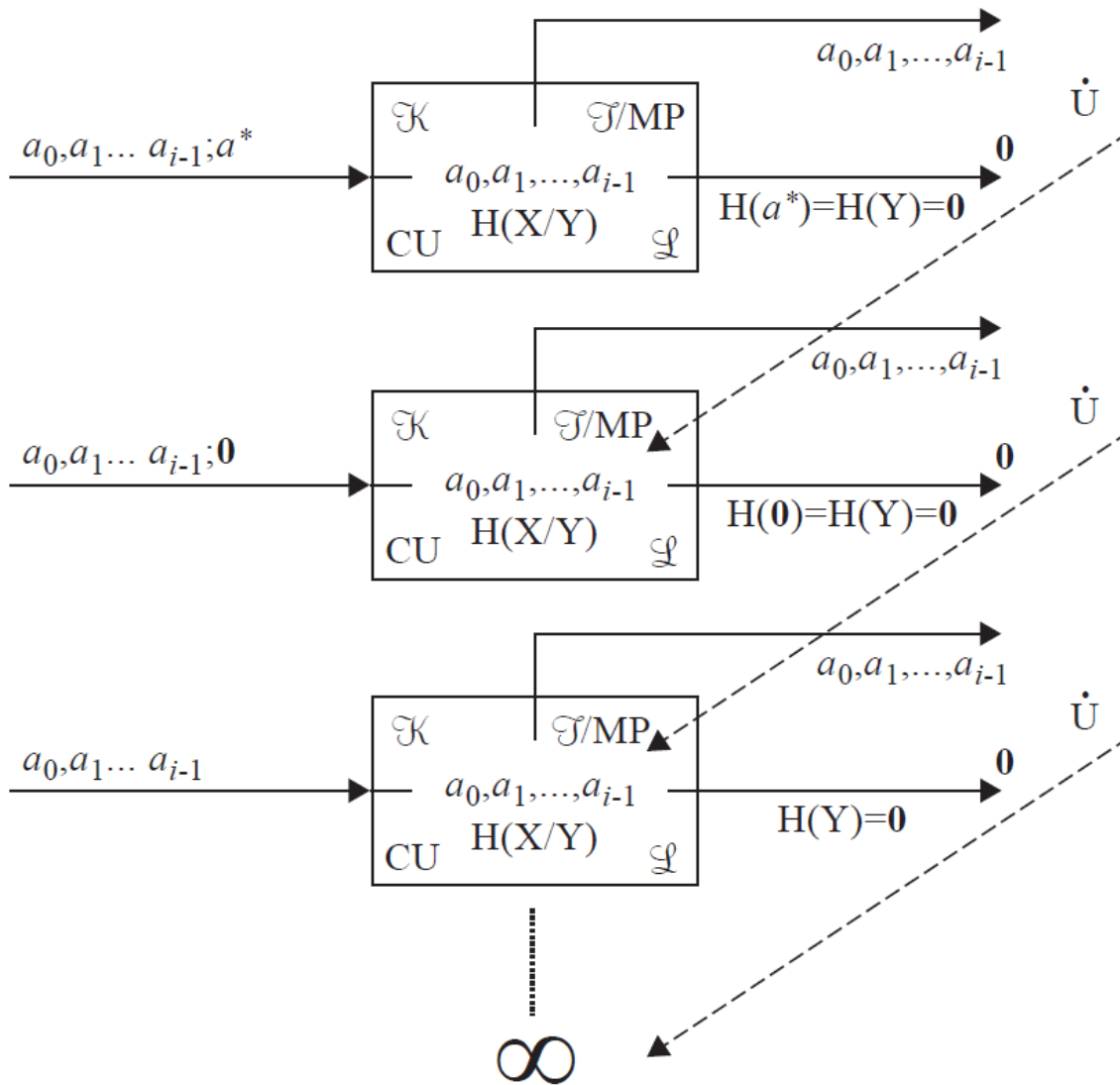
any \mathcal{K} can't transfer its own states²¹

or observe/copy/measure itself.

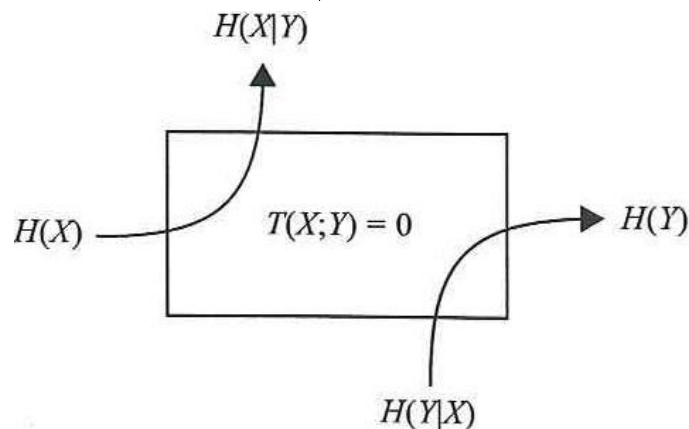
²⁰It is against the Caratheodory theorems. The existence of the *Perpetuum Mobile II.* and *I.* is required. It also requires the time arrow change $\frac{S_c}{t} > 0$. 'Solving' this 'problem' represents the belief in the Maxwell demon's functionality. The need of distinguishing between the measured and the measuring - to avoid the *HALTING PROBLEM* - leads to the formulation of the Gödel theorems and their physical form as the Caratheodory theorems and vice versa.

²¹Used as input messages. *There's a need for a 'step-aside' outside the measured to gain a not zero and positive result* $T(X;Y) = H(X) - H(X|Y) = H(Y) > 0$

The last *CLAIM* a^* in the input $\vec{x} = \overrightarrow{a_0, a_1, \dots, a_{i-1}; a^*}$ is not inferrable and, as such, interrupts the channel \mathcal{K} ,

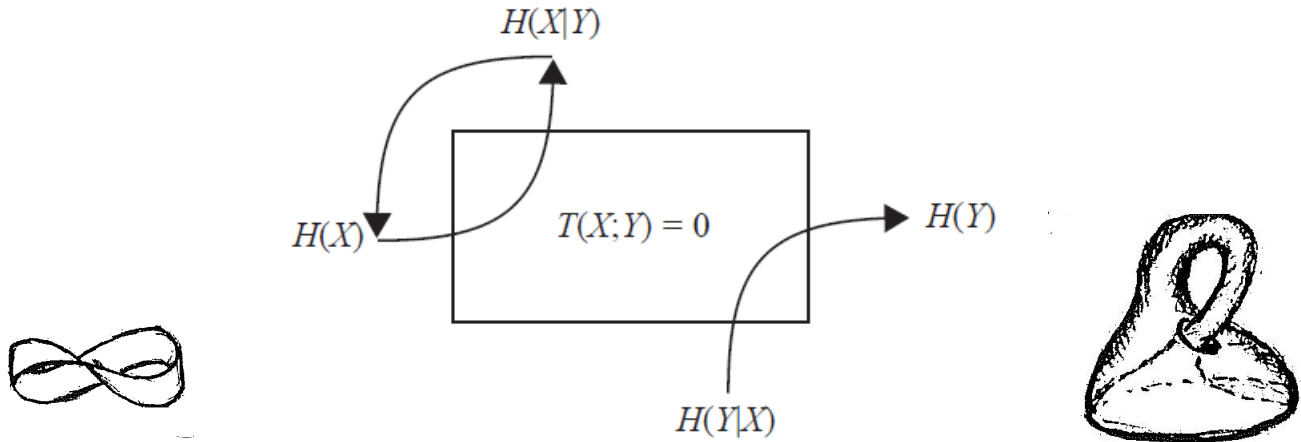


$\mathbf{H(X)}$ expresses the \mathcal{K} 's/ TM 's structure $\mathbf{H(X|Y)}$,



$\mathbf{H(Y)}$ does not relate to $\mathbf{H(X)}$

Repeating this **autoreference** attempt
causes the **infinite cycle**,²²



$H(X) = H(X|Y)$ is the **HALTING PROBLEM** core.

Our bad awaiting
of a not zero and positive output, $H(Y) > 0$,
formulates the methodological error, but,
identifiable by the II. P.T.

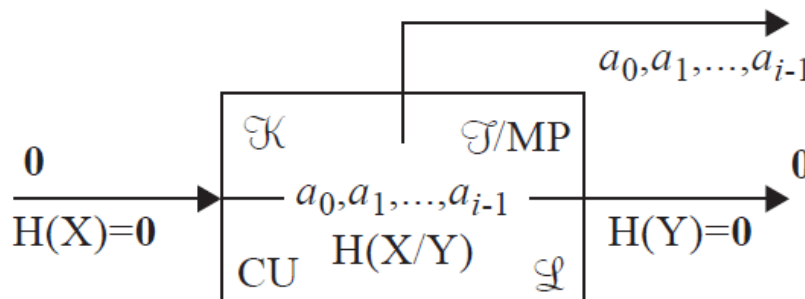
by this possible consideration,

$$H(X|Y) \cong H(a_0, a_1, \dots, a_{i-1}) > 0, \quad \underline{H(X) = 0}$$

$$H_c = \underline{H(X)} - H(X|Y) = \underline{0} - H(a_0, a_1, \dots, a_{i-1}) < 0$$

which **violates/contradicts** - $\Delta S_c < 0$ - the **II. P.T.**

while we have



²²The same $H(Y) = H(Y|X)$, $H(X|Y)$ and the whole configuration (type) are repeated infinitely.

11. Autoreference and Gödel

Now let \mathbf{x} be the *SEQUENCE OF FORMULAE* **valid ("1")** in the theory $\mathcal{T}_{\mathcal{P}\mathcal{A}}$ and \mathbf{y} is a *CLAIM* ²³ **with a general syntax of $\mathcal{P}^\star - \mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$, $\mathcal{P}^\star \supset \mathcal{P} \supset \mathcal{T}_{\mathcal{P}\mathcal{A}}$.**²⁴

We define the **valid ("1")** relation $\mathbf{Q}(\mathbf{x}, \mathbf{Y})$ saying:

" $\mathbf{x} \in \mathbb{X}$ is not in the $\mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ -INFERENCE relation to the values \mathbf{y} of \mathbf{Y} ", $[\mathbf{y}|\mathbf{Y} \equiv \mathcal{T}_{\mathcal{P}\mathcal{A}}\text{-property}]$

$\mathbf{Y} : \mathbb{Y} \cong \{\{\Omega_{\mathcal{E}^\star}\}^\star - \{l_{\Omega_{\mathcal{E}}}\}\} \ni \mathbf{y}, \quad \mathbf{X} : \mathbb{X} \stackrel{\Delta}{=} \mathcal{T}_{\mathcal{P}\mathcal{A}} \cong \{l_{\Omega_{\mathcal{E}}}\} \ni \mathbf{x}$

$\mathbf{Q}(\mathbf{x}, \mathbf{Y}) = \text{"1"}$, from the construction and, **supposedly,**

$\mathbf{Q}(\mathbf{x}, \mathbf{y}) = \text{"1"}$ \equiv **Proof $_{\mathcal{P}}$ [$\mathbf{Q}(\mathbf{x}, \mathbf{y})$]** for \mathbf{x} and any $\mathbf{y} \in \mathbf{Y}$

But, when we set

$\underline{\mathbf{y}} = \mathbf{Q}(\mathbf{x}, \mathbf{Y}) = \mathbf{y}(\mathbf{Y})$ $[\cong \mathbf{y}|\mathbf{Y} \equiv \mathcal{T}_{\mathcal{P}\mathcal{A}}\text{-property}]$

$\mathbf{Q}(\mathbf{x}, \mathbf{y}) = \mathbf{y}[\mathbf{Q}(\mathbf{x}, \mathbf{Y})] = \mathbf{y}(\mathbf{y}) = \mathbf{Q}[\mathbf{x}, \mathbf{Q}(\mathbf{x}, \mathbf{Y})]$

the $\mathbf{Y} := \mathbf{y}$ in $\mathbf{y}(\mathbf{Y})$ generates the *autoreference*,

$\mathbf{Q}(\mathbf{x}, \mathbf{y})$; $\mathbf{Q}[\mathbf{x}, [\mathbf{Q}(\mathbf{x}, \mathbf{Y})]]$, $\mathbf{Q}[\mathbf{x}, [\mathbf{Q}(\mathbf{x}, [\mathbf{Q}(\mathbf{x}, \mathbf{Y})])]]$, ...

neither constructible in \mathcal{P} and nor provable by $\mathcal{T}_{\mathcal{P}\mathcal{A}}$ but its *validity follows* from the *II. P.T. validity*.²⁵

$\overline{\text{Proof}_{\mathcal{P}}[\mathbf{Q}(\mathbf{x}, \mathbf{y})]} = \text{"1"}$

II.P.T. $\cong [\text{Proof}_{\mathcal{P}}[\mathbf{Q}(\mathbf{x}, \mathbf{y})] = \text{"0"}] = \text{"1"}$

$[[\text{Proof}_{\mathcal{E}}([\mathbf{d}]\mathbf{Q}_{\text{Ext}, l_{\Omega_{\mathcal{E}}}} = \text{"0"}) = \text{"0"}] = \text{"1"}] \equiv \text{II. P.T.}$

²³The \mathbf{x} and \mathbf{y} are the Gödel numbers.

²⁴ $\mathcal{P}^\star \cong \{\Omega_{\mathcal{E}^\star}\}^\star, \mathcal{P} \cong \{\Omega_{\mathcal{E}}\}, \mathcal{T}_{\mathcal{P}\mathcal{A}} \cong \{l_{\Omega_{\mathcal{E}}}\}, x \cong l_{\Omega_{\mathcal{E}}}$.

²⁵And within the PM - *Principia Mathematica: B. Russel, L. Whitehead.*

Now we set

$$\underline{\mathbf{p}} = [\forall_{\mathbf{x} \in \mathbb{X}} | \mathbf{Q}(\mathbf{x}, \mathbf{Y})] \stackrel{\Delta}{=} \mathbf{Q}(\mathbb{X}, \mathbf{Y}) = \mathbf{p}(\mathbf{Y})$$

"no $\mathbf{x} \in \mathbb{X}$ is in the $\mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ -INFERENCE relation to the variable \mathbf{Y} " - to its space $\mathbb{Y} \cong \{ \{\Omega_{\mathcal{E}^*}\}^{\star} - \{\mathbf{l}_{\Omega_{\mathcal{E}}}\} \}$

With the substitution $\mathbf{Y} := \mathbf{p}$

the resulting $\mathbf{Q}(\mathbb{X}, \mathbf{p}) = \mathbf{p}(\mathbf{p})$ [$\cong \mathbf{p} | \mathbf{Y} \equiv \mathcal{T}_{\mathcal{P}\mathcal{A}}$ -property]

contains the *autoreference*

- the *totality* $\mathbb{X}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ 'defines' (??) its property \mathbf{Y} -

$$[\forall_{\mathbf{x} \in \mathbb{X}} | [\mathbf{Q}[\mathbf{x}, [\forall_{\mathbf{x} \in \mathbb{X}} | \mathbf{Q}(\mathbf{x}, \mathbf{Y})]]]]_{\mathbf{Y} := \mathbf{p}} = \mathbf{Q}[\mathbb{X}, [\mathbf{Q}(\mathbb{X}, \mathbf{Y})]]_{\mathbf{Y} := \mathbf{p}}$$

$$\underline{\mathbf{Q}(\mathbb{X}, \mathbf{p})}; \quad \mathbf{Q}[\mathbb{X}, [\mathbf{Q}(\mathbb{X}, \mathbf{Y})]], \quad \mathbf{Q}[\mathbb{X}, [\mathbf{Q}(\mathbb{X}, [\mathbf{Q}(\mathbb{X}, \mathbf{Y})])]], \quad \dots$$

As before, by the *II. P.T.*'s validity, we know:

$$\overline{\mathbf{Proof}_{\mathcal{P}}[\mathbf{Q}(\mathbb{X}, \mathbf{p})]} = \text{"1"}$$

$$\text{II.P.T.} \cong [\mathbf{Proof}_{\mathcal{P}}[\mathbf{Q}(\mathbb{X}, \mathbf{p})] = \text{"0"}] = \text{"1"}$$

$$[[\mathbf{Proof}_{\mathcal{E}}([\mathbf{d}] \mathbf{Q}_{\text{Ext}, \{\mathbf{l}_{\Omega_{\mathcal{E}}}\}} = \mathbf{0}) = \text{"0"}] = \text{"1"}] \equiv \text{II. P.T.}$$

$$\mathbf{y}, \mathbf{p}, \mathbf{Q}(\mathbf{x}, \mathbf{y}), \mathbf{Q}(\mathbb{X}, \mathbf{p}) \in \{\mathcal{P}^{\star} - \mathcal{P}\} \cong \{ \{\Omega_{\mathcal{E}^*}\}^{\star} - \Omega_{\mathcal{E}} \},$$

$$\{\mathbf{l}_{\Omega_{\mathcal{E}}}\} \cong \mathcal{T}_{\mathcal{P}\mathcal{A}} \stackrel{\Delta}{=} \mathbb{X}, \quad \underline{\mathbf{Q}(\mathcal{T}_{\mathcal{P}\mathcal{A}}, \mathbf{p}) = \mathbf{Q}(\mathbb{X}, \mathbf{p}) \stackrel{\Delta}{=} \mathbf{17Gen} \mathbf{r} \cong \mathbf{Y}},$$

$$\mathbf{r} \stackrel{\Delta}{=} \mathbf{r}(\mathbf{X}) = \mathbf{Q}(\mathbf{X}, \mathbf{p}), \quad \mathbf{17Gen} \mathbf{r} \stackrel{\Delta}{=} \forall_{\mathbf{x} \in \mathbb{X}} | \mathbf{r}(\mathbf{X}), \quad \text{card } \mathbf{Y} = \aleph_1$$

The consistent theory $\mathcal{T}_{\mathcal{P}_A}$ is constructed

by ourselves, from its outside,

it does **not contain the axiom of its consistency**

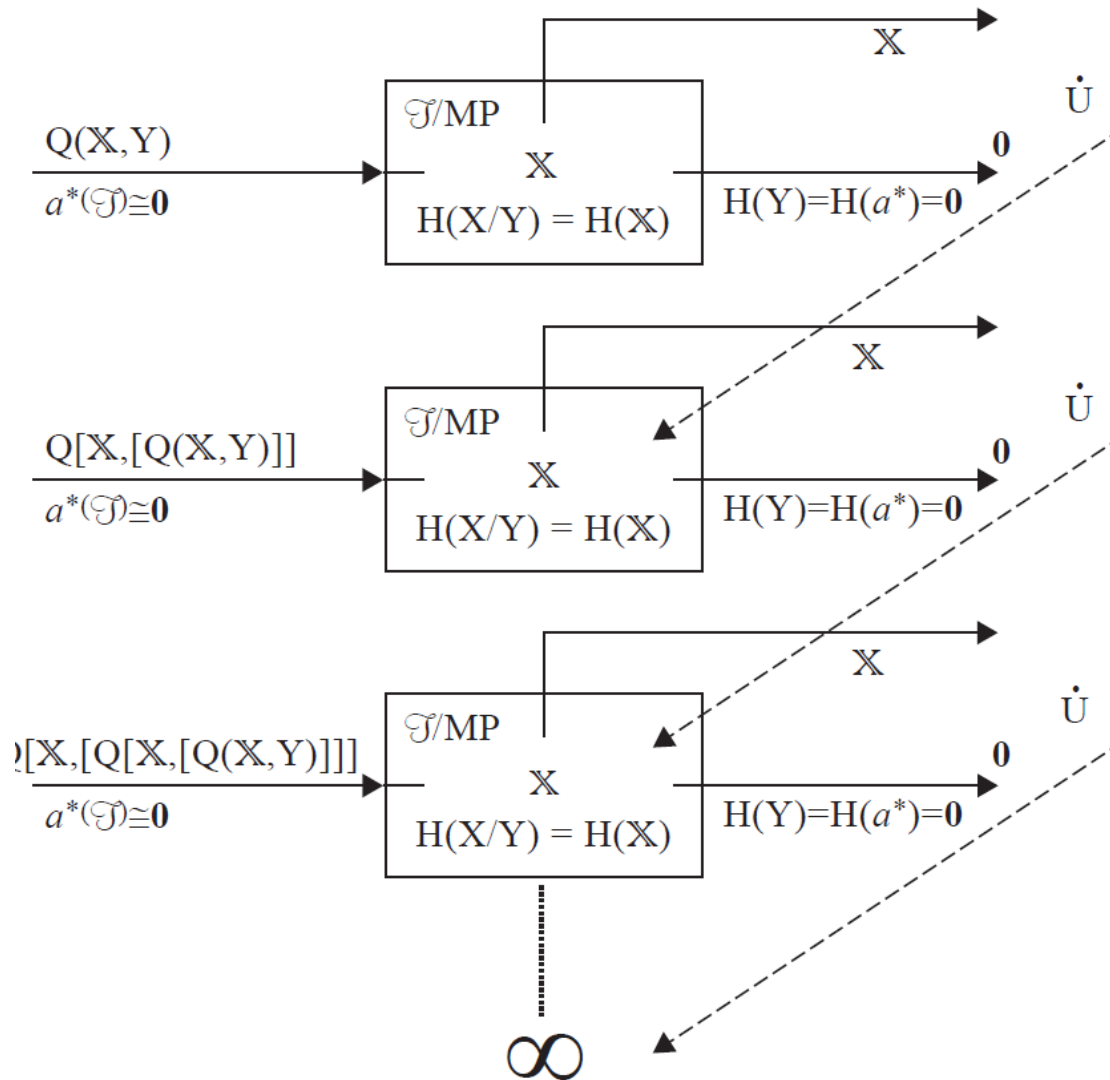
as its own object/formula/status

The attempt to prove the $\mathcal{T}_{\mathcal{P}_A}$ consistency by it itself

- imbuilding is property (Y) to itself as its own object -

defines the autoreference/*HALTING PROBLEM*,

not resultative but, as such, *identifiable*;



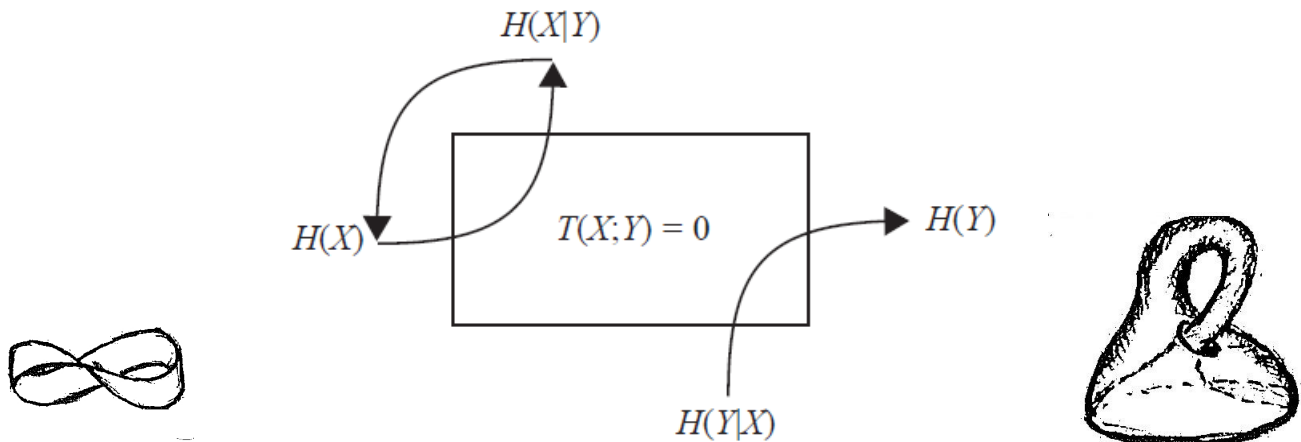
The *CLAIM* $Q(\mathbb{X}, p)/17Gen r$ is constructed purely syntactically as for $\mathcal{P}/\mathcal{T}_{\mathcal{P}_A}$, in $\{\mathcal{P}^* - \mathcal{P}\}$,

but is the descriptive *CLAIM*

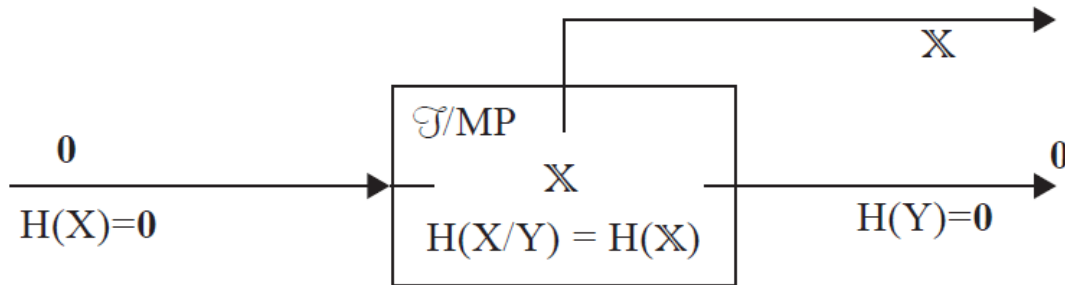
of the $\mathcal{T}_{\mathcal{P}_A} [\mathbb{X}]$ consistent inference in \mathcal{P} .²⁶

²⁶Also the Rao-Cramer inequality from the Mathematical Statistics documents it.

For the channel \mathcal{K} , associated to the inferring TM , is being interrupted infinitely,²⁷ in any transfer act, the TM and the \mathcal{K} realize the Klein bottle run.



But, in fact we have



We formulate the methodological error for \mathcal{T}_{PA} but identifiable by the II. P.T.'s violation, $\Delta S_c < 0$,

$$\underline{H(X) = H[Q(X, Y)] = 0, \quad H(X|Y) \cong H(X/\mathcal{T}_{PA}) > 0}$$

$$H_c = H(X) - H(X|Y) = \underline{0} - H(X/\mathcal{T}_{PA}) < 0$$

The another form of the II. P.T. is formulated:

$$Q(X, p) / \underline{17Gen r = "1"}^{28}$$

²⁷It is the reason for the non expressibility of the consistency of the theory \mathcal{T}_{PA} CLAIM - of the axiom of its consistency expressing its general property - in it itself. The theory \mathcal{T}_{PA} does not contain this axiom as its own formula.

²⁸The same methodological error represent the Epimenides liar and the Richard paradox - the property of a certain totality is not formulable within it as its own object part.

12. Appendix 1 - Gödel Theorems

I. Gödel theorem (corrected semantically by [19] and [23]):²⁹

For every *recursive* and *consistent*

CLASS OF FORMULAE κ ,³⁰

and outside this set³¹,

exists the true ("1") **CLAIM** \mathbf{r}

$[\mathbf{r} \notin \kappa, \mathbf{r} \notin \mathbf{Flg}(\kappa), \mathbf{r} \in \mathbf{Y}, \text{card } \mathbf{Y} = \aleph_1]$

with a free **VARIABLE** \mathbf{v} that

neither the **CLAIM** $\mathbf{vGen } \mathbf{r}$

nor the **CLAIM** $\mathbf{Neg}(\mathbf{vGen } \mathbf{r})$

belongs to the set $\mathbf{Flg}(\kappa)$

$[\mathbf{vGen } \mathbf{r} \notin \mathbf{Flg}(\kappa)] \ \& \ [\mathbf{Neg}(\mathbf{vGen } \mathbf{r}) \notin \mathbf{Flg}(\kappa)],$

$\mathbf{r} \triangleq \mathbf{r}(\mathbf{X}) = \mathbf{Q}[\mathbf{X}, \forall_{\mathbf{x} \in \mathbb{X}} | \mathbf{Q}(\mathbf{x}, \mathbf{Y})], \quad \mathbf{17Gen } \mathbf{r} \triangleq \forall_{\mathbf{x} \in \mathbb{X}} | \mathbf{r}(\mathbf{X})$

CLAIMS $\mathbf{vGen } \mathbf{r}$ and $\mathbf{Neg}(\mathbf{vGen } \mathbf{r})$

are not κ -**PROVABLE**

the **CLAIM** $\mathbf{vGen } \mathbf{r}$ *is not* κ -**DECIDABLE**.

[They are elements of the formulating/syntactic
meta-system κ^\star , inconsistent against κ .]

²⁹Gödel-Rosser theorem.

³⁰Recursively axiomatizable and with the given set of the inference rules (Peano/Robinson arithmetics). $\kappa = \mathcal{P}$ is from the PM - *Principia Mathematica: B. Russel, L. Whitehead*. 1910, 1912, 1913, 1927. , containing the Peano arithmetics $\mathcal{T}_{\mathcal{P}\mathcal{A}}$. On the $\mathcal{T}_{\mathcal{P}\mathcal{A}}$ the *real and complex number arithmetics* is based. The PM is the textit formal-syntax-logic-semantic base for the physical theories/hypotheses.

³¹Far from "...[PA-]arithmetic and sentencial/*SENTENCIAL*" and far from (!) "In"

II. Gödel theorem (corrected semantically by [19] and [23]):³²

If κ is an arbitrary *recursive and consistent*

CLASS OF FORMULAE,

then any **CLAIM** y [$y \in Y$, $\text{card } Y = \aleph_1$]

saying that **CLASS** κ is consistent

must be constructed outside this set³³

and for this fact, it is not κ -**PROVABLE**³⁴

The consistency of the **CLASS OF FORMULAE** κ
is tested by the *relation* Wid(κ)³⁵

$$\text{Wid}(\kappa) \sim (\mathbf{Ex})[\text{CLAIM}(y) \ \& \ \overline{\text{Proof}_\kappa(y)}]$$

The **FORMULAE** class κ is consistent

\Leftrightarrow

at least one κ -**UNPROVABLE CLAIM** y exists.

Now $y = \underline{17\text{Gen } r} \notin \mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$, $\kappa = \mathcal{T}_{\mathcal{P}\mathcal{A}} \subset \mathcal{P} \subset \mathcal{P}^\star$

$$[[\text{Proof}_{\mathcal{P}}(\underline{17\text{Gen } r}) = "0"] = "1"] \equiv \text{Wid}(\mathcal{T}_{\mathcal{P}\mathcal{A}})$$

$$\cong$$

$$[[\text{Proof}_{\mathcal{G}}([\text{d}]\text{Q}_{\text{Ext}, \Omega_{\mathcal{G}}} = 0) = "0"] = "1"] \equiv \text{II. P.T.}$$

³²Gödel-Rosser theorem.

³³Far from "...[PA-]arithmetic and sentential/*SENTENCIAL*" and far from (!) "In"

³⁴Any attempt to prove/*PROVE/INFER* it in the system \mathcal{P}/κ leads to the requirement for inconsistency of the consistent (!) system \mathcal{P}/κ (in fact we are entering into the inconsistent meta-system \mathcal{P}^\star - see the real sense of the Proposition V.).

³⁵Die *Widerspruchsfreiheit* - the Consistency.

13. Appendix 2

Under the adiabaticity, $[d]Q_{\text{Ext}}=0$, of the system \mathcal{L}
it is not possible to derive such a *CLAIM*
that is stating this adiabatic supposition.

This *CLAIM* is constructible *not adiabatically*,
outside the adiabatic \mathcal{L} only.

Autoreference/*HALTING PROBLEM*

Self-Observation

- the *CLAIM* about adiabaticity of \mathcal{L} *within* \mathcal{L} -
- the *CLAIM* about consistency of $\mathcal{T}_{\mathcal{P}\mathcal{A}}$ *within* \mathcal{P} -
is excluded.

This is the *nature law* expressed
by the *Caratheodory formulation* of the *II. P.T.*
and by the *Gödel theorems' sense*.

The eye can not look at and into itself.

Any mixing of the various
observation/expressing/approach levels
leads to the paradoxes and is to be excluded.

Under the consistency of the system \mathcal{P}
it is not possible to derive such a *CLAIM*
that is stating this consistency supposition.

This *CLAIM* is constructible *purely syntactically*,
outside the consistent \mathcal{P} only (in $\mathcal{P}^\star - \mathcal{P}$).³⁶

³⁶The Great Fermat's theorem is not inferrable within $\mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ either - it is not of the $\mathcal{P}/\mathcal{T}_{\mathcal{P}\mathcal{A}}$ type, although is arithmetical. In fact it is not a part of the \mathcal{P} but of the $\mathcal{P}^\star - \mathcal{P}$, $\mathcal{P}^\star \cong \{\mathbf{p}, \mathbf{V}, \mathbf{T}\} = \mathbb{R}^3$.



(1).PDF

That's me or it is the picture of me - P1.



(2).PDF

This is the mirror picture of me - P2.



(3).PDF

Here I have 'ordered' the mirror picture P2 to step out from the mirror a stand, e.g., in front of me/P1 having the right hands overlapped.

CHAOS, EQUILIBRIUM, INFINITE CYCLE, PARADOX by mixing of various observation levels.

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